

**Some Graph, Digraph, and Mixed  
Graph Results Concerning  
Decompositions, Packings, and  
Coverings**

(ABSTRACT #1003-05-120)

Robert Gardner

Benedict Bobga and Chau Nguyen

Department of Mathematics

East Tennessee State University

Gary Coker

Department of Computer Science

Francis Marion College

presented at

Joint Meeting of the

A.M.S. and the M.A.A.

*AMS Special Session on*

*Design Theory and Graph Theory, I*

Atlanta, Georgia

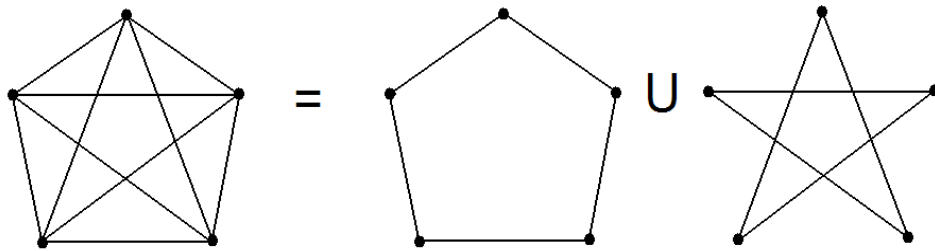
January 5, 2005

# 1. Decompositions

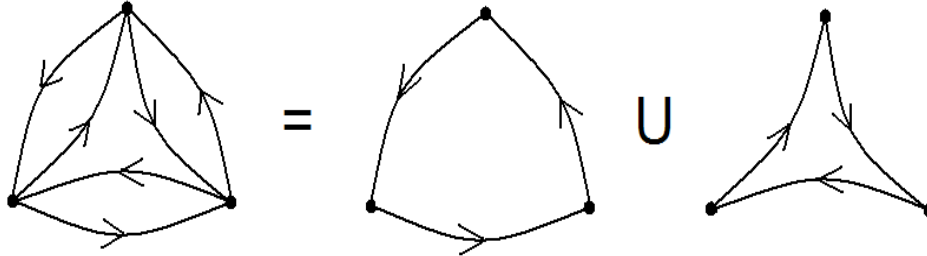
**Definition.** A *decomposition* of a simple graph  $G$  with isomorphic copies of graph  $g$  is a set  $\{g_1, g_2, \dots, g_n\}$  where  $g_i \cong g$  and  $V(g_i) \subset V(G)$  for all  $i$ ,  $E(g_i) \cap E(g_j) = \emptyset$  if  $i \neq j$ , and  $\cup_{i=1}^n g_i = G$ . Here,  $V(G)$  is the vertex set of graph  $G$  and  $E(G)$  is the edge set of graph  $G$ .

**Note.** Decompositions of digraphs are similarly defined (replacing edge sets with arc sets).

**Example.** There is a decomposition of  $K_5$  into 5-cycles ( $C_5$ 's):

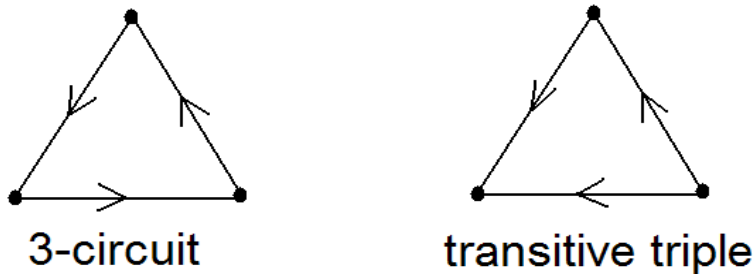


**Example.** There is a decomposition of the complete digraph on 3 vertices ( $D_3$ ) into 3-circuits:



**Note.** A decomposition of  $K_v$  into  $C_3$ 's is equivalent to a *Steiner triple system* of order  $v$ , denoted  $STS(v)$ . It is well known that a  $STS(v)$  exists if and only if  $v \equiv 1$  or  $3 \pmod{6}$ .

**Definition.** There are two orientations of  $C_3$ , a 3-circuit and a transitive triple:



A decomposition of  $D_v$  into 3-circuits is equivalent to a *Mendelsohn triple system* of order  $v$ ,  $MTS(v)$ . A decomposition of  $D_v$  into transitive triples is equivalent to a *directed triple system* of order  $v$ ,  $DTS(v)$ . A  $MTS(v)$  exists if and only if  $v \equiv 0$  or  $1 \pmod{3}$ ,  $v \neq 6$  [Mendelsohn, 1971]. A  $DTS(v)$  exists if and only if  $v \equiv 0$  or  $1 \pmod{3}$  [Hung and Mendelsohn, 1973].

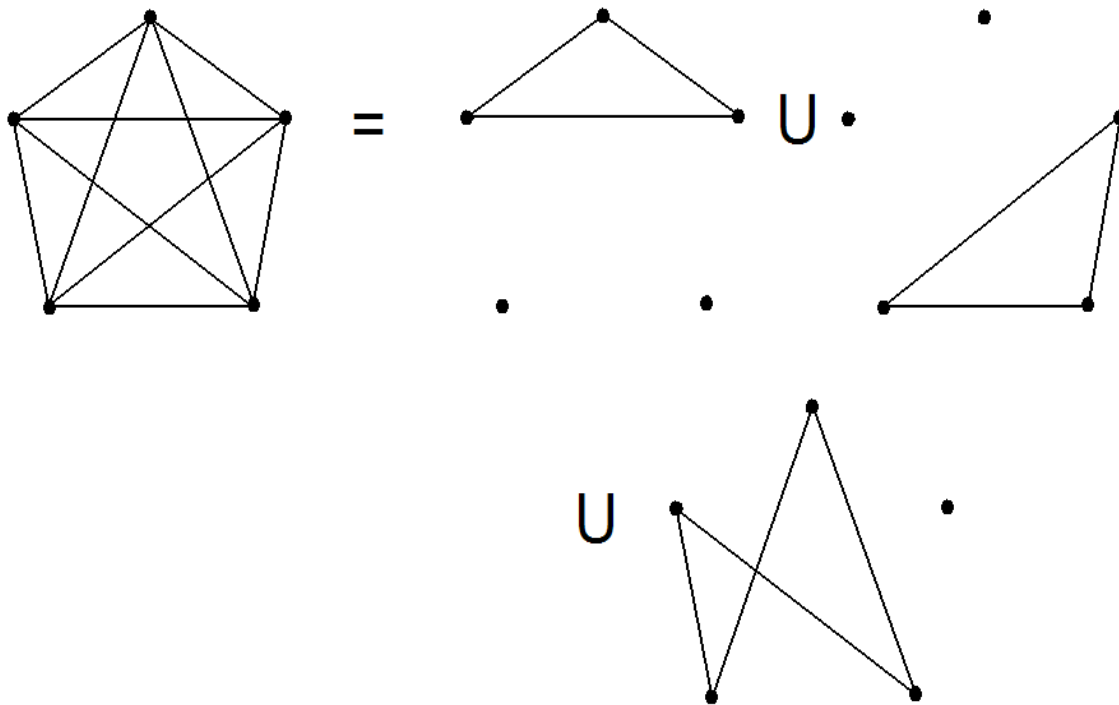
## 2. Packings and Coverings

**Definition.** A *maximal packing* of a simple graph  $G$  with isomorphic copies of graph  $g$  is a set  $\{g_1, g_2, \dots, g_n\}$  where  $g_i \cong g$  and  $V(g_i) \subset V(G)$  for all  $i$ ,  $E(g_i) \cap E(g_j) = \emptyset$  if  $i \neq j$ ,  $\cup_{i=1}^n g_i \subset G$ , and

$$|E(L)| = |E(G) \setminus \cup_{i=1}^n E(g_i)|$$

is minimal. The set  $L$  is called the *leave* of the packing.

**Example.** A packing of  $K_5$  with 3-cycles has a leave  $L$  with 4 edges:



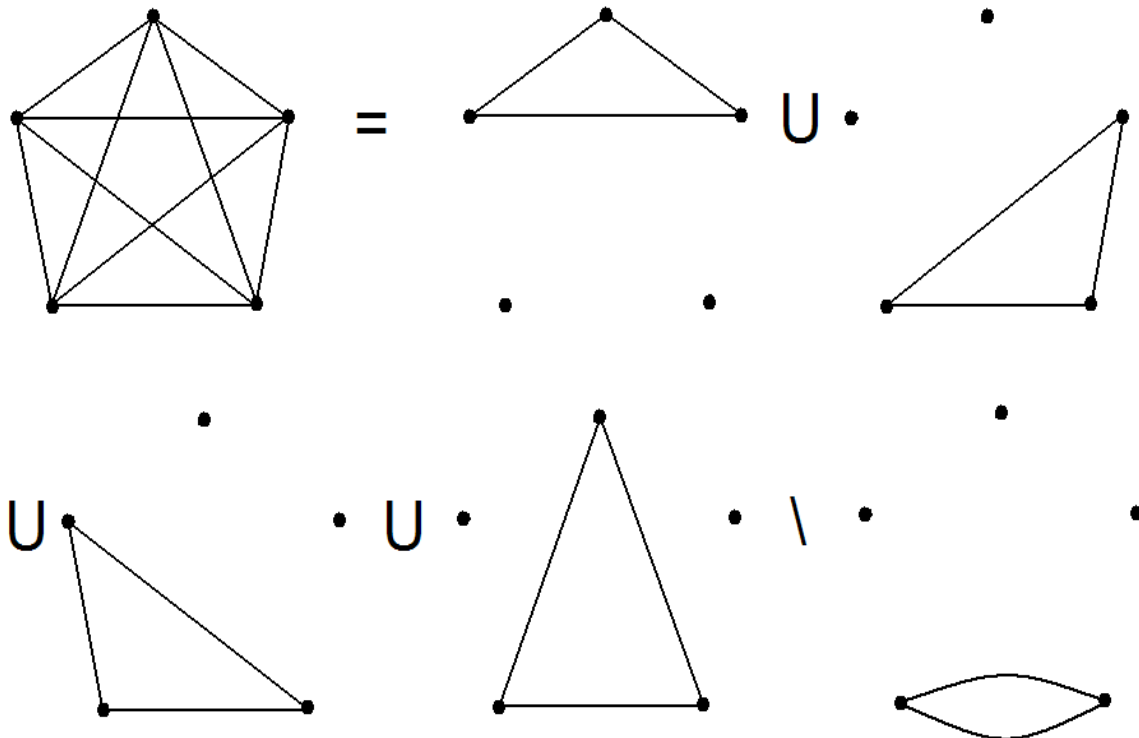
**Note.** Packings of the complete graph on  $v$  vertices,  $K_v$ , with graph  $g$  have been studied for  $g$  a 3-cycle [Schönheim, 1966],  $g$  a 4-cycle [Schönheim and Bialostocki, 1975],  $g = K_4$  [Brouwer, 1979], and  $g$  a 6-cycle [Kennedy, 1993].

**Definition.** A *minimal covering* of a simple graph  $G$  with isomorphic copies of a graph  $g$  is a set  $\{g_1, g_2, \dots, g_n\}$  where  $g_i \cong g$  and  $V(g_i) \subset V(G)$  for all  $i$ ,  $G \subset \cup_{i=1}^n g_i$ , and

$$|E(P)| = |\cup_{i=1}^n E(g_i) \setminus E(G)|$$

is minimal (the graph  $\cup_{i=1}^n g_i$  may not be simple and  $\cup_{i=1}^n E(g_i)$  may be a multiset). The graph  $P$  is called the *padding* of the covering.

**Example.** A covering of  $K_5$  with 3-cycles has a padding of  $2 \times K_2$ :



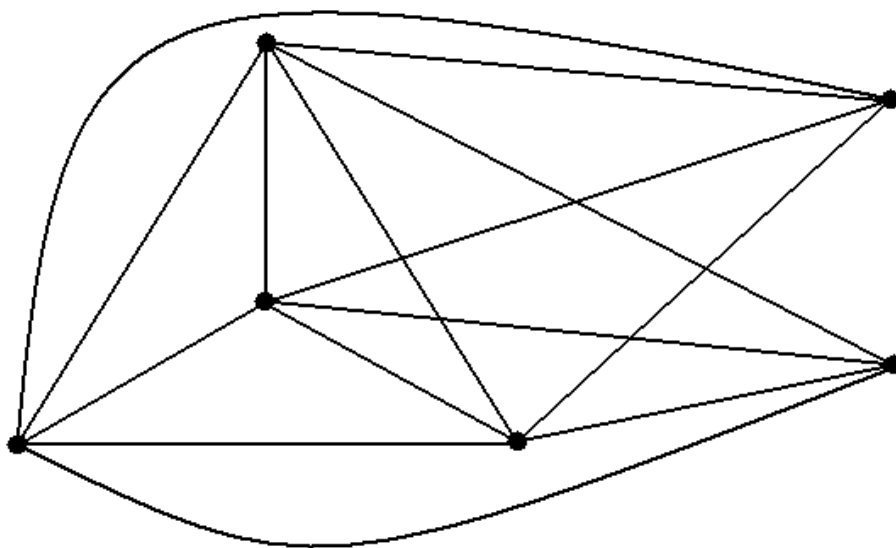
**Note.** Coverings of  $K_v$  with graph  $g$  have been studied for  $g$  a 3-cycle [Fort and Hedlund, 1958],  $g$  a 4-cycle [Schönheim and Bialostocki, 1975], and  $g$  a 6-cycle [Kennedy, 1995].

### 3. 4-Cycles and the Complete Graph with a Hole

**Definition.** The *complete graph on  $v$  vertices with a hole of size  $w$* , denoted  $K(v, w)$  is the graph with vertex set  $V(K(v, w)) = V_{v-w} \cup V_w$  where  $|V_{v-w}| = v - w$  and  $|V_w| = w$ , and with edge set  $E(K(v, w)) = \{(a, b) \mid a \neq b, \{a, b\} \subset V_{v-w} \cup V_w \text{ and } \{a, b\} \not\subset V_w\}$ .

We let  $V_{v-w} = \{1_1, 2_1, \dots, (v-w)_1\}$  and  $V_w = \{1_2, 2_2, \dots, w_2\}$ .

**Example.** The complete graph on 6 vertices with a hole of size 2,  $K(6, 2)$ , is:



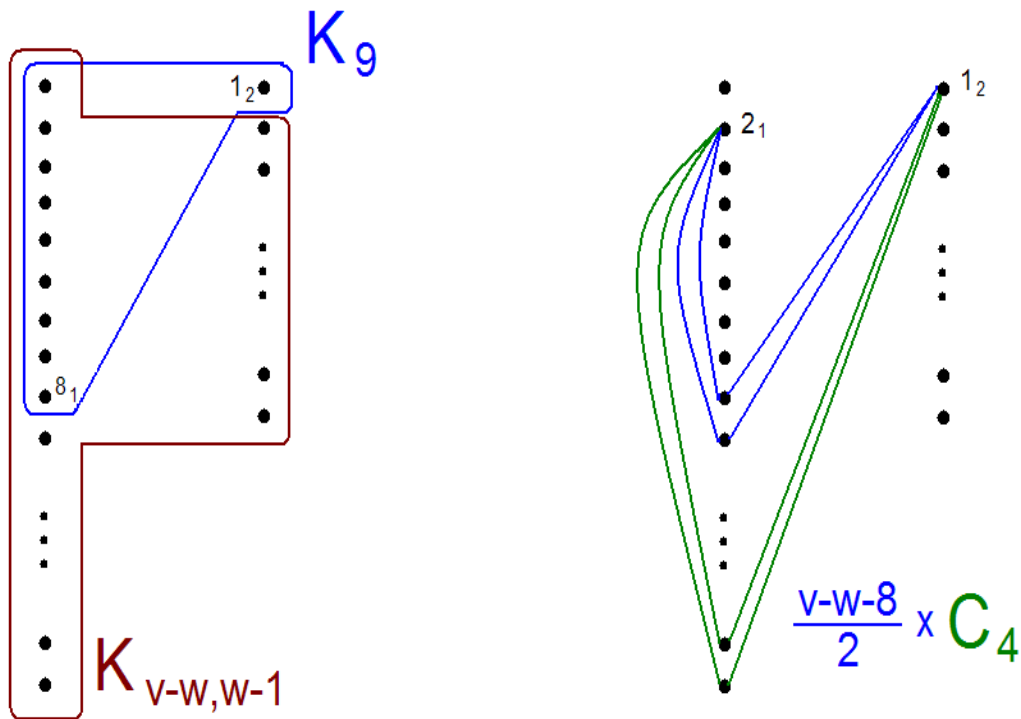
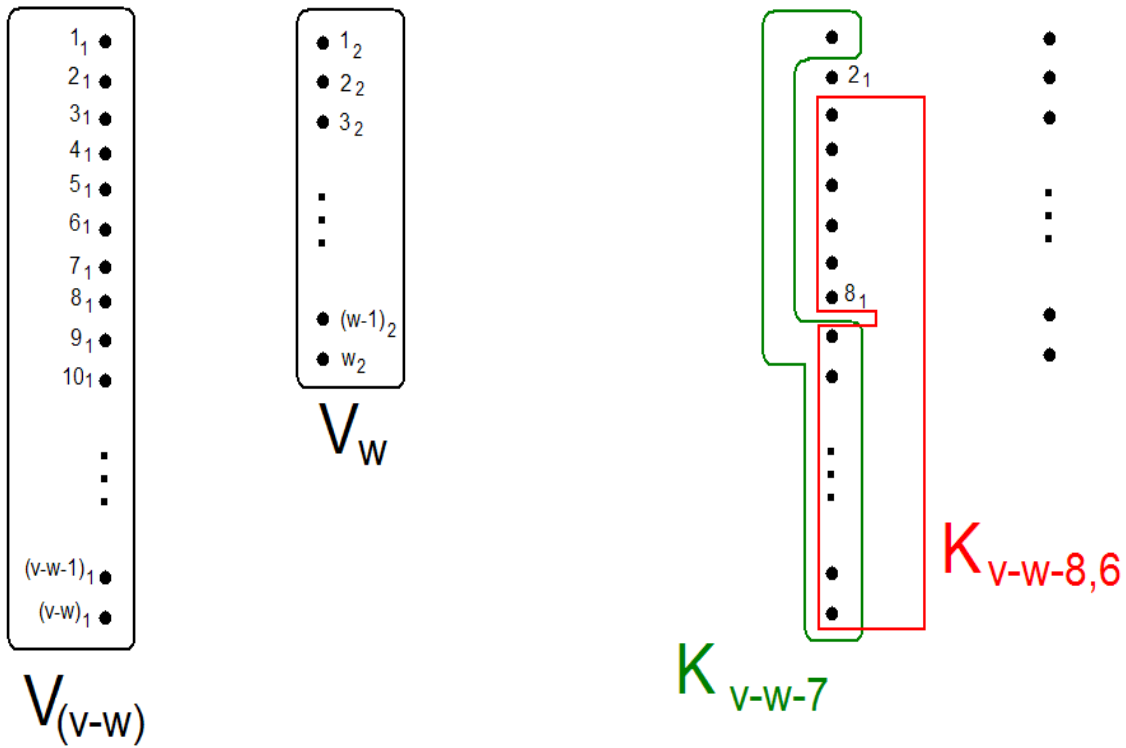
**Note.** It is rather well known that a 4-cycle decomposition of  $K_v$  exists if and only if  $v \equiv 1 \pmod{8}$  [Schönheim and Bialostocki, 1975]. It is quite easy to show that  $K_{m,n}$  can be decomposed into  $C_4$ 's if and only if  $m \equiv n \equiv 0 \pmod{2}$ .

**Theorem.** [Gardner, Lavoie, and Nguyen, 2005] A  $C_4$  decomposition of  $K(v, w)$  exists if and only if  $w \equiv 1 \pmod{2}$  and  $v - w \equiv 0 \pmod{8}$ .

**Proof.** Since each vertex of  $C_4$  is even, a necessary condition is that each vertex of  $K(v, w)$  must be even. A vertex of  $V_{v-w}$  is of degree  $v - 1$ , therefore  $v$  must be odd. A vertex of  $V_w$  is of degree  $v - w$  and so  $v - w$  must be even and  $w$  must be odd. The graph  $K(v, w)$  has  $v(v - 1)/2 - w(w - 1)/2$  edges. Since  $C_4$  has four edges, another necessary condition for the desired decomposition is that  $v(v - 1)/2 - w(w - 1)/2 \equiv 0 \pmod{4}$ . Together, these conditions yield the necessary conditions of the theorem.

Now

$$K(v, w) = K_{v-w-7} \cup K_{v-w-8,6} \cup K_{v-w,w-1} \cup K_9 \cup (v-w-8)/2 \times C_4.$$



Now  $v-w-7 \equiv 1 \pmod{8}$  and  $9 \equiv 1 \pmod{8}$ , so  $K_{v-w-7}$  and  $K_9$  can be decomposed into  $C_4$ 's. Next,  $v-w$  and  $w-1$  are even,  $v-w-8$  and  $6$  are even, so  $K_{v-w,w-1}$  and  $K_{v-w-8,6}$  can be decomposed into  $C_4$ 's. Therefore  $K(v, w)$  can be decomposed into  $C_4$ 's. ■



**Note.** Schönheim and Bialostocki [1975] studied  $C_4$  packings of  $K_v$ . Bryant and Khodkar [2000] studied  $C_3$  packings of  $K(v, w)$ . We now look at  $C_4$  packings of  $K(v, w)$ .

**Theorem.** [Gardner, Lavoie, and Nguyen, 2005] A  $C_4$  packing of  $K(v, w)$  exists if and only if:

1. if  $v - w \equiv 0 \pmod{2}$  and  $w \equiv 1 \pmod{2}$ , then

$$|E(L)| = \begin{cases} 0 & \text{if } v - w \equiv 0 \pmod{8} \\ 3 & \text{if } v - w \equiv 2 \pmod{8} \\ 6 & \text{if } v - w \equiv 4 \pmod{8} \\ 5 & \text{if } v - w \equiv 6 \pmod{8}, \end{cases}$$

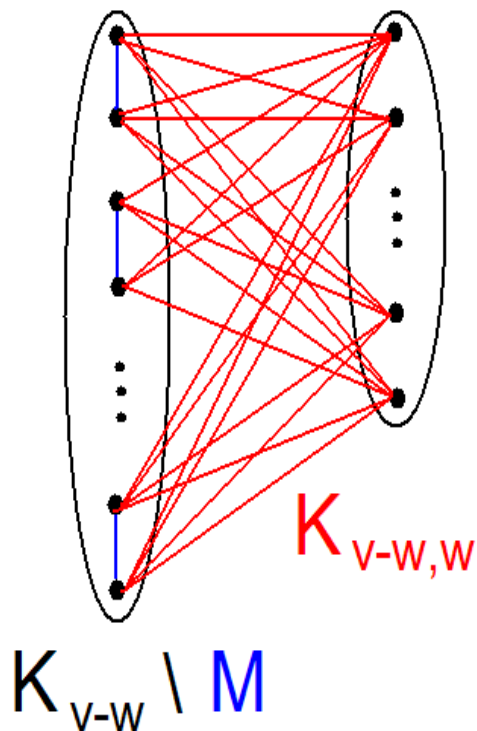
2. if  $v - w \equiv 0 \pmod{2}$  and  $w \equiv 0 \pmod{2}$ , then  $|E(L)| = (v - w)/2$ ,

3. if  $v - w \equiv 1 \pmod{2}$  and  $w \equiv 0 \pmod{2}$ , then  $|E(L)| = w + k$  where  $k$  is the minimum nonnegative integer such that  $|E(K(v, w))| - |E(L)| \equiv 0 \pmod{4}$ ,

4. if  $v - w \equiv 1 \pmod{2}$ ,  $w \equiv 1 \pmod{2}$ , and  $v - w \leq w$ , then  $|E(L)| = w + k$  where  $k$  is the minimum nonnegative integer such that  $|E(K(v, w))| - |E(L)| \equiv 0 \pmod{4}$ , and

5. if  $v - w \equiv 1$ ,  $w \equiv 1 \pmod{2}$ , and  $v - w > w$ , then  $|E(L)| = v/2 + k$  where  $k$  is the minimum nonnegative integer such that  $|E(K(v, w))| - |E(L)| \equiv 0 \pmod{4}$ .

**“Proof.”** The proof consists of 17 cases. Consider the case  $v - w \equiv 0 \pmod{2}$  and  $w \equiv 0 \pmod{2}$ . Each vertex of  $V_{v-w}$  is of degree  $v - w - 1$  which is odd, therefore in the leave  $L$  each vertex from  $V_{v-w}$  must be of odd degree. So a packing with  $|E(L)| = (v - w)/2$  would be optimal. Now  $K(v, w) = (K_{v-w} \setminus M) \cup K_{v-w,w} \cup M$ :



We can show (using difference methods) that  $K_{v-w} \setminus M$  can be decomposed into  $C_4$ 's. Since  $v - w$  and  $w$  are both even, then  $K_{v-w,w}$  can be decomposed into  $C_4$ 's. We then have an optimal packing with leave  $L = M$  where  $M$  is a matching on  $K_{v-w}$  and so  $|E(L)| = (v - w)/2$ . ■

**Theorem.** [Gardner, Lavoie, and Nguyen, 2005] A  $C_4$  covering of  $K(v, w)$  exists if and only if:

1. if  $v - w \equiv 0 \pmod{2}$ ,  $v - w > 2$ , and  $w \equiv 1 \pmod{2}$ , then

$$|E(P)| = \begin{cases} 0 & \text{if } v - w \equiv 0 \pmod{8} \\ 5 & \text{if } v - w \equiv 2 \pmod{8} \\ 2 & \text{if } v - w \equiv 4 \pmod{8} \\ 3 & \text{if } v - w \equiv 6 \pmod{8}, \end{cases}$$

2. if  $v - w \equiv 0 \pmod{4}$  and  $w \equiv 0 \pmod{2}$ , then  $|E(P)| = (v - w)/2$ ,

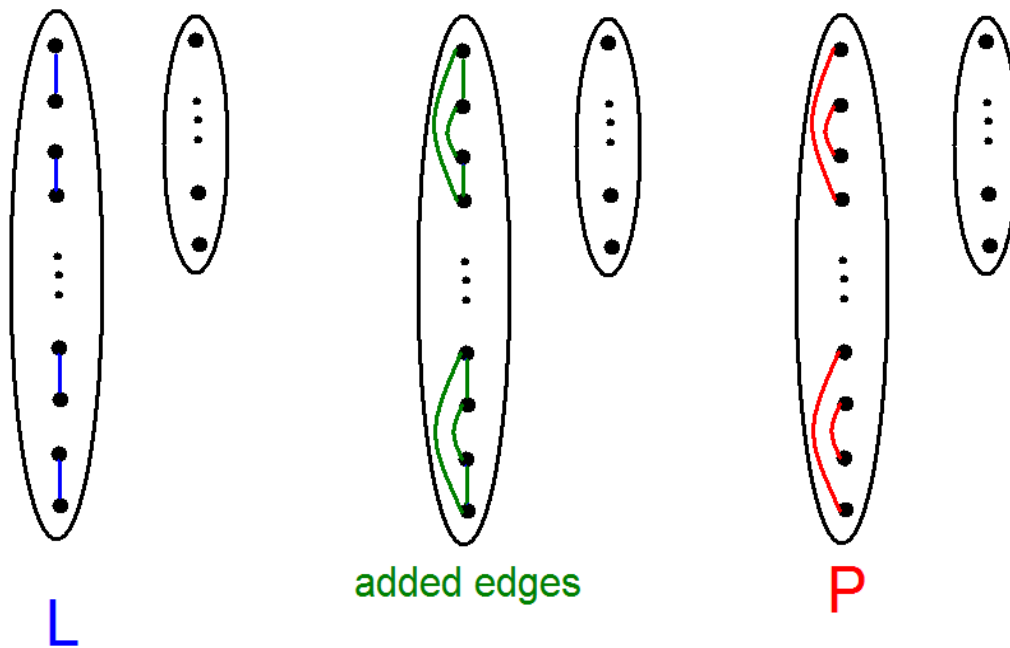
3. if  $v - w \equiv 2 \pmod{4}$  and  $w \equiv 0 \pmod{2}$ , then  $|E(P)| = (v - w)/2 + 2$ ,

4. if  $v - w \equiv 1 \pmod{2}$ ,  $v - w > 1$ , and  $w \equiv 0 \pmod{2}$ , then  $|E(P)| = w + k$  where  $k$  is the minimum nonnegative integer such that  $|E(K(v, w))| + |E(P)| \equiv 0 \pmod{4}$ ,

5. if  $v - w \equiv 1 \pmod{2}$ ,  $v - w > 1$ ,  $w \equiv 1 \pmod{2}$ , and  $v - w \leq w$ , then  $|E(P)| = w + k$  where  $k$  is the minimum nonnegative integer such that  $|E(K(v, w))| + |E(P)| \equiv 0 \pmod{4}$ , and

6. if  $v - w \equiv 1$ ,  $w \equiv 1 \pmod{2}$ , and  $v - w > w$ , then  $|E(P)| = v/2 + k$  where  $k$  is the minimum nonnegative integer such that  $|E(K(v, w))| + |E(P)| \equiv 0 \pmod{4}$ .

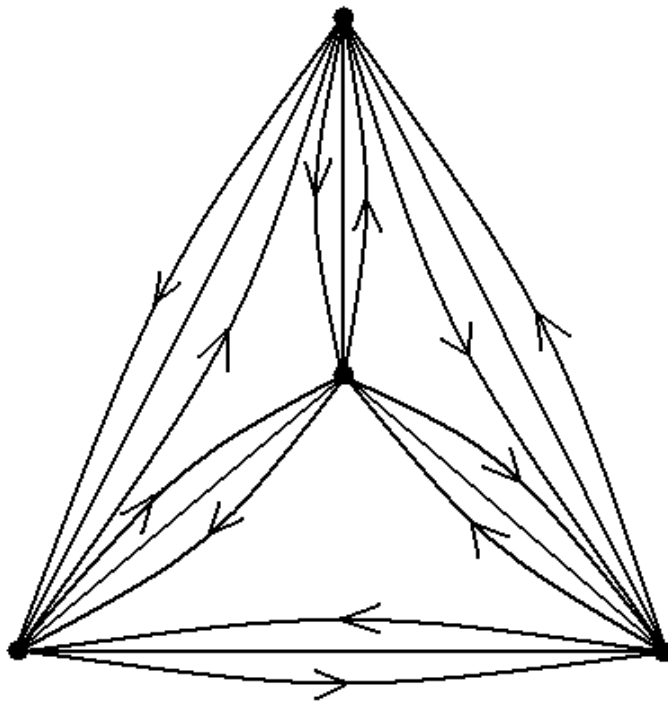
“**Proof.**” The proof consists of 22 cases. Consider the case  $v - w \equiv 0 \pmod{2}$  and  $w \equiv 0 \pmod{2}$ . Each vertex of  $V_{v-w}$  is of degree  $v - w - 1$  which is odd, therefore in the padding  $P$  each vertex from  $V_{v-w}$  must be of odd degree. So a covering with  $|E(P)| = (v - w)/2$  would be optimal. We take the packing described above with the leave  $L$  a matching on  $V_{v-w}$ . We then add  $C_4$ 's as below, and have a padding  $P$  which is also a matching on  $V_{v-w}$  and hence  $|E(P)| = (v - w)/2$ . ■



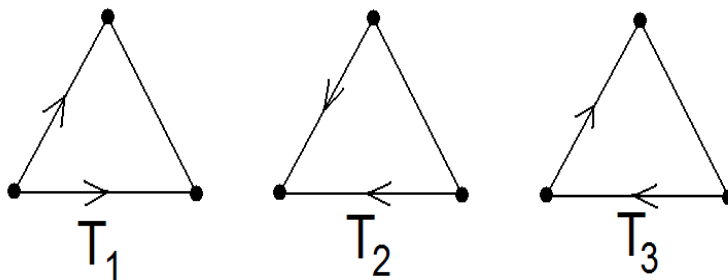
## 4. Some Results Concerning Mixed Graphs

**Definition.** A *mixed graph* on  $v$  vertices is an ordered pair  $(V, C)$  where  $V$  is a set of vertices,  $|V| = v$ , and  $C$  is a set of ordered and unordered pairs, denoted  $[x, y]$  and  $(x, y)$  respectively, of elements of  $V$ . An ordered pair  $[x, y] \in C$  is called an *arc* of  $(V, C)$  and an unordered pair  $(x, y) \in C$  is called an *edge* of graph  $(V, C)$ . The *complete mixed graph* on  $v$  vertices, denoted  $M_v$ , is the mixed graph  $(V, C)$  where, for every pair of distinct vertices  $v_1, v_2 \in V$ , we have  $\{[v_1, v_2], [v_2, v_1], (v_1, v_2)\} \subset C$ .

**Example.** The mixed graph  $M_4$  is:



**Note.** Since  $M_v$  has twice as many arcs as edges, we are inspired to study triple system based on complete mixed graphs and the following:



**Definition.** A decomposition of  $M_v$  into  $T_i$ 's is a  $T_i$  triple system or order  $v$ .

**Theorem.** [Gardner, 1999] A  $T_i$  triple system of order  $v$  exists for all  $i \in \{1, 2, 3\}$  and  $v \equiv 1 \pmod{2}$ , except for  $i = 3$  and  $v \in \{3, 5\}$ .

**Note.** A study of packing and covering  $M_v$  with  $T_i$  ( $i \in \{1, 2, 3\}$ ) is currently underway by Bobga and Gardner.

**Definition.** Let  $G$  be a graph and  $\gamma = \{g_1, g_2, \dots, g_n\}$  be a  $g$  decomposition of  $G$ . An *automorphism* of this decomposition is a permutation of  $V(G)$  which fixes set  $\gamma$ . An automorphism of digraph and mixed graph decompositions are similarly defined.

**Definition.** Consider a permutation on a set of size  $v$ . The permutation is said to be *cyclic* if it consists of a single cycle of length  $v$ . It is *bicyclic* if it consists of two disjoint cycles of lengths  $N_1$  and  $N_2$  where  $v = N_1 + N_2$ .

**Theorems.** A cyclic  $STS(v)$  exists if and only if  $v \equiv 1$  or  $3 \pmod{6}$ ,  $v \neq 9$  [Peltesohn, 1939]. A bicyclic  $STS(v)$  where  $v = N_1 + N_2$  admitting an automorphism whose disjoint cyclic decomposition is a cycle of length  $N_1$ , where  $N_1 > 1$ , and a cycle of length  $N_2$  exists if and only if  $N_1 \equiv 1$  or  $3 \pmod{6}$ ,  $N_1 \neq 9$ ,  $N_1 \mid N_2$ , and  $v = N_1 + N_2 \equiv 1$  or  $3 \pmod{6}$  [Calahan-Zijlstra and Gardner, 1994].

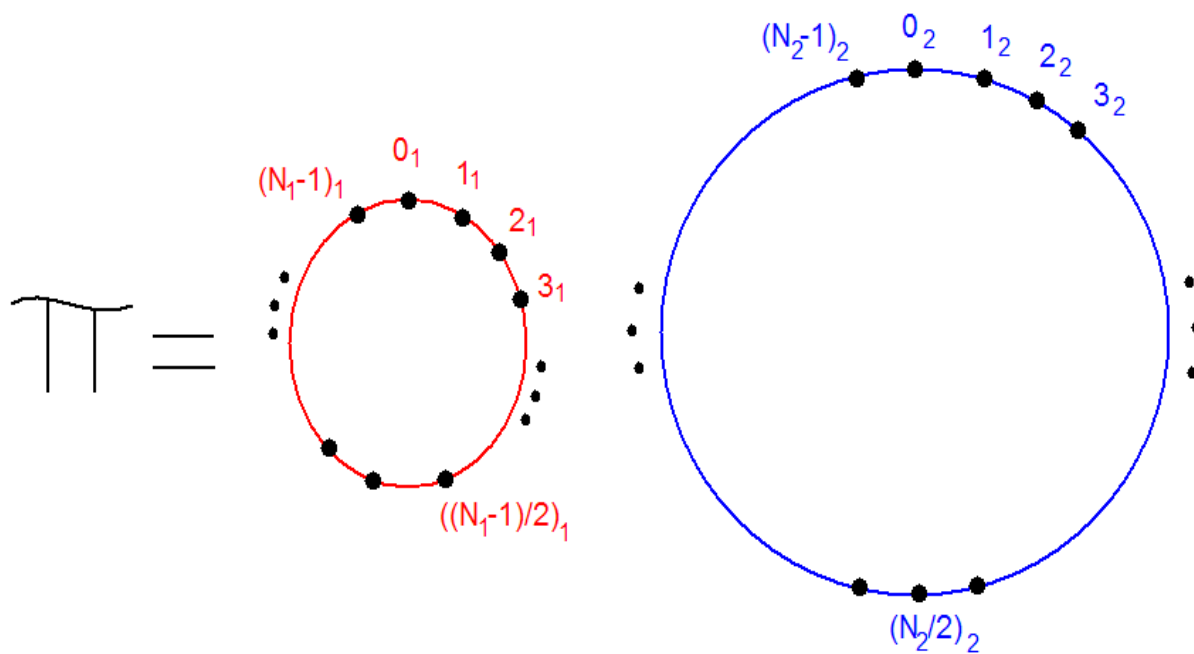
**Theorems.** A cyclic  $DTS(v)$  exists if and only if  $v \equiv 1, 4,$  or  $7 \pmod{12}$  [Colbourn and Colbourn, 1982]. A bicyclic  $DTS(v)$  admitting an automorphism consisting of two cycles each of length  $v/2$  exists if and only if  $v \equiv 4 \pmod{6}$ . A bicyclic  $DTS(v)$  admitting an automorphism consisting of a cycle of length  $N_1$  and a cycle of length  $N_2$ , where  $v = N_1 + N_2$ , exists if and only if  $N_1 \equiv 1, 4,$  or  $7 \pmod{12}$  and  $N_2 = kN_1$  where  $k \equiv 2 \pmod{3}$  [Gardner, 1998].

**Theorem.** A cyclic  $MTS(v)$  exists if and only if  $v \equiv 1$  or  $3 \pmod{6}$ ,  $v \neq 9$  [Colbourn and Colbourn, 1981]. To the best of my knowledge, bicyclic  $MTS$ 's have not been studied.

**Theorem.** [Gardner, 1999] A cyclic  $T_i$  triple system of order  $v$  exists for all  $i \in \{1, 2, 3\}$  and  $v \equiv 1 \pmod{2}$ , except for  $i = 3$  and  $v \in \{3, 5\}$ .

**Theorem.** [Bobga and Gardner, 2005] A bicyclic  $T_i$  triple system, where  $i \in \{1, 2\}$ , exists admitting an automorphism consisting of a cycle of length  $N_1$  and a cycle of length  $N_2$ , where  $N_1 < N_2$ , if and only if  $N_1 \equiv 1 \pmod{2}$ ,  $N_1 \mid N_2$ , and  $v = N_1 + N_2 \equiv 1 \pmod{2}$ . A bicyclic  $T_3$  triple system does not exist.

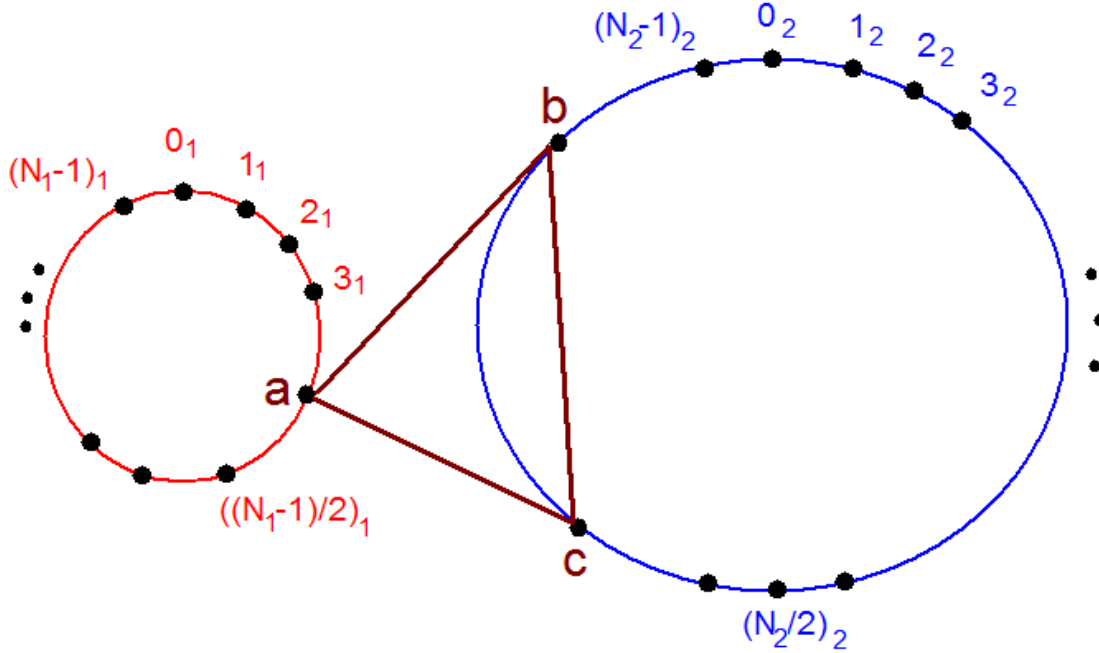
**“Proof.”** Let  $\pi$  be a bicyclic automorphism of a  $T_i$  system where  $\pi$  consists of disjoint cycles of lengths  $N_1$  and  $N_2$ :



Notice that  $\pi^{N_1}$  fixes the points  $\{0_1, 1_1, \dots, (N_1 - 1)_1\}$ . Therefore these points form a cyclic subsystem of order  $N_1$  and hence  $N_1 \equiv 1 \pmod{2}$ .



Now consider some  $T_i$  with vertex set  $\{a, b, c\}$  where  $a \in \{0_1, 1_1, \dots, (N_1 - 1)_1\}$  and  $\{b, c\} \subset \{0_2, 1_2, \dots, (N_2 - 1)_2\}$ :



When we apply  $\pi^{N_2}$  to this triple, we see that  $\pi^{N_2}(b) = b$  and  $\pi^{N_2}(c) = c$ , and hence  $\pi^{N_2}(a) = a$ . That is,  $N_2$  is a multiple of  $N_1$ . This established the necessary conditions. Sufficiency is established through difference methods. ■

## 5. References

1. B. Bobga and R. Gardner, Bicyclic, Rotational, and Reverse Mixed Triple Systems, in preparation.
2. D. Bryant and A. Khodkar, Maximum Packings of  $K_v - K_u$  with Triples, *Ars Combinatoria* **55** (2000), 259–270.
3. A. Brouwer, Optimal Packings of  $K_4$ 's into a  $K_n$ , *Journal of Combinatorial Theory, Series A* **26**(3) (1979), 278–297.
4. R. Calahan-Zijlstra, Bicyclic Steiner Triple Systems, *Discrete Math.* **128** (1994), 35–44.
5. M. Colbourn and C. Colbourn, Disjoint Cyclic Mendelsohn Triple Systems, *Ars Combinatoria* **11** (1981), 3–8.
6. M. Colbourn and C. Colbourn, The Analysis of Directed Triple Systems by Refinement, *Annals of Discrete Math.* **15** (1982), 97–103.
7. M. Fort and G. Hedlund, Minimal Coverings of Pairs by Triples, *Pacific Journal of Mathematics* **8** (1958), 709–719.
8. R. Gardner, Bicyclic Directed Triple Systems, *Ars Combinatoria*, **49** (1998) 249–257.
9. R. Gardner, Triple Systems from Mixed Graphs, *Bulletin of the ICA* **27** (1999), 95–100.
10. R. Gardner, C. Nguyen, S. Lavoie, 4-Cycle Packings and Coverings of the Complete Graph with a Hole, in preparation.
11. S.H.Y. Hung and N.S. Mendelsohn, Directed Triple Systems, *Journal Combinatorial Theory, Series A* **14** (1973), 310–318.
12. J. Kennedy, Maximum Packings of  $K_n$  with Hexagons, *Australasian Journal of Combinatorics* **7** (1993), 101–110.
13. J. Kennedy, Two Perfect Maximum Packings and Minimum Coverings of  $K_n$  with Hexagons, Ph.D. dissertation, Auburn University, U.S.A. 1995.
14. N.S. Mendelsohn, A Natural Generalization of Steiner Triple Systems, “Computers in Number Theory,” eds. A.O. Atkins and B. Birch, Academic Press, London, 1971.
15. R. Pelsesohn, Eine Lösung der beiden Heffterschen Differenzenprobleme, *Compositio Math.* **6** (1939), 251–257.
16. J. Schönheim, On Maximal Systems of  $k$ -Tuples, *Studia Sci. Math. Hungarica* (1966), 363–368.
17. J. Schönheim and A. Bialostocki, Packing and Covering of the Complete Graph with 4-Cycles, *Canadian Mathematics Bulletin* **18**(5) (1975), 703–708.