## Graph Theory

#### Chapter 12. Stable Sets and Cliques 12.3. Ramsey's Theorem—Proofs of Theorems











**Theorem 12.9.** For any two integers  $k \ge 2$  and  $\ell \ge 2$ ,

$$r(k,\ell) \leq r(k,\ell-1) + r(k-1,\ell).$$

Furthermore, if  $r(k, \ell - 1)$  and  $f(k - 1, \ell)$  are both even, strict inequality holds in the inequality.

**Proof.** Let G be a graph on  $r(k, \ell - 1) + r(k - 1, \ell)$  vertices and let  $v \in V$ . We consider two cases:

1. Vertex v is nonadjacent to a set S of at least  $r(k, \ell - 1)$  vertices.

2. Vertex v is adjacent to a set T of at least  $r(k-1, \ell)$  vertices.

Since G has  $r(k, \ell - 1) + r(k - 1, \ell)$  vertices, then there are  $r(k, \ell - 1) + r(k - 1, \ell) - 1$  vertices in G other than vertex v. So the number of vertices to which v is nonadjacent plus the number of vertices to which v is adjacent is equal to  $r(k, \ell - 1) + r(k - 1, \ell) - 1$ ; hence, either Case 1 or Case 2 must hold.

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# Theorem 12.9 (continued 1)

**Proof (continued).** In Case 1, the induced subgraph G[S] of G contains either a clique of k vertices or a stable set of  $\ell - 1$  vertices. Therefore,  $G[S \cup \{v\}]$  contains either a clique of k vertices (since G[S] does) or a stable set of  $\ell$  vertices (the stable set of  $\ell - 1$  vertices in G[S] along with vertex v which is not adjacent to any vertices of S). Since  $G[S \cup \{v\}]$  is a subgraph of G, then G also contains these sets. In Case 2, the induced subgraph G[T] contains either a clique of k-1 vertices or a stable set of  $\ell$  vertices. Therefore,  $G[T \cup \{v\}]$  contains either a clique of k vertices (the clique of k - 1 vertices in G[T] along with vertex v which is adjacent to all vertices of T) or a stable set of  $\ell$  vertices (since G[T] does). Since  $G[T \cup \{v\}]$  is a subgraph of G, then G also contains these sets. This proves the inequality.

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# Theorem 12.9 (continued 2)

**Proof (continued).** Now suppose that  $r(k, \ell - 1)$  and  $r(k - 1, \ell)$  are both even, and let G' be a graph on  $r(k, \ell - 1) + r(k - 1, \ell) - 1$  vertices. So G' has an odd number vertices, by Corollary 1.2 there is some vertex v'of G' of even degree. In particular, v' cannot be adjacent to precisely  $r(k-1, \ell) - 1$  vertices. So v' must be adjacent to at least vertices  $r(k-1, \ell)$  vertices (in which Case 2 above holds) or v' must be nonadjacent to at least  $r(k, \ell - 1)$  vertices (in which Case 1 above holds). That is, in graph G' either Case 1 or Case 2 hold, and hence, as shown above, G' either contains a clique on k vertices or a stable set on  $\ell$ vertices. So we have by the inequality established above (but with  $r(k-1,\ell)$  there replaced with  $r(k-1,\ell)-1$  here) gives

$$r(k,\ell) \leq r(k,\ell-1) + r(k-1,\ell) - 1 < r(k,\ell-1) + r(k-1,\ell),$$

as claimed.

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as claimed.

**Theorem 12.10.** For all positive integers k and  $\ell$ ,  $r(k, \ell) \leq \binom{k+\ell-2}{k-1}$ .

**Proof.** We give an inductive proof on the sum  $k + \ell$ . If  $k + \ell \le 5$  then either k or  $\ell$  must be less than 3. By Note 12.A, for  $k + \ell \le 5$  we have  $r(1,\ell) = 1 \le \binom{k+\ell-2}{k-1}$  since every combination is at least 1,  $r(k,2) = k = \binom{k+\ell-2}{k-1} = \binom{k}{k-1} = k$ , and similarly, by symmetry, r(k,1) and  $r(2,\ell)$  are also bounded as claimed. So we take  $k + \ell \le 5$  as the base case(s). **Theorem 12.10.** For all positive integers k and  $\ell$ ,  $r(k, \ell) \leq \binom{k+\ell-2}{k-1}$ .

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Let *m* and *n* be positive integers and for the induction hypothesis suppose the theorem is valid for all integers *k* and  $\ell$  such that  $5 \le k + \ell \le m + n$ .

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# Theorem 12.10 (continued)

**Theorem 12.10.** For all positive integers k and  $\ell$ ,  $r(k, \ell) \leq \binom{k + \ell - 2}{k - 1}$ . **Proof (continued).** Then

$$\begin{aligned} r(m,n) &\leq r(m,n-1) + r(m-1,n) \text{ by Theorem 12.9} \\ &\leq \binom{m+n-3}{m-1} + \binom{m+n-3}{m-2} \text{ by the induction hypothesis} \\ &= \frac{(m+n-3)!}{(n-2)!(m-1)!} + \frac{(m+n-3)!}{(n-1)!(m-2)!} \\ &= \frac{(m+n-3)!((n-1)+(m-1))}{(n-1)!(m-1)!} = \frac{(m+n-2)!}{(n-1)!(m-1)!} \\ &= \binom{m+n-2}{m-1}. \end{aligned}$$

So the induction step is established. Therefore, by mathematical induction, the claim holds for all positive integers k and  $\ell$ .

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### **Theorem 12.12.** For all positive integers k, $r(k, k) \ge 2^{k/2}$ .

**Proof.** By Note 12.A, r(1,1) = 1 and r(2,2) = 2, so we just need to consider  $k \ge 3$ . Let  $\mathcal{G}_n$  be the set of all simple graphs with vertex set  $\{v_1, v_2, \ldots, v_n\}$ . Let  $\mathcal{G}_n^k$  be the set of these labeled simple graphs which have a clique on k vertices. We have  $\mathcal{G}_n = 2^{\binom{n}{2}}$  (since for any of the  $\binom{n}{2}$  pairs of vertices may or may not be joined by an edge; see Note 1.2.B).

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$$1 \cdot 2^{\binom{n-k}{2}} \cdot 2^{(n-k)k} = 2^{\binom{n^2-n-k^2+k}{2}} = 2^{\binom{n}{2} - \binom{k}{2}}$$

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# Theorem 12.12 (continued 1)

**Theorem 12.12.** For all positive integers k,  $r(k, k) \ge 2^{k/2}$ .

**Proof (continued).** Because there are  $\binom{n}{k}$  distinct *k*-element subsets of  $\{v_1, v_2, \ldots, v_n\}$  we have  $|\mathcal{G}_n^k| \leq \binom{n}{k} 2^{\binom{n}{2} - \binom{k}{2}}$ . We pick up an inequality here because there may be graphs in  $\mathcal{G}_n^k$  which have more than one *k*-clique in which case they are counted once on the left-hand-side by the bound on the right-hand-side counts then more once. Therefore

$$\begin{aligned} \frac{|\mathcal{G}_n^k|}{|\mathcal{G}_n|} &\leq \frac{\binom{n}{k} 2\binom{n}{2} - \binom{k}{2}}{2\binom{n}{2}} = \binom{n}{k} 2^{-\binom{k}{2}} \\ &= \frac{n(n-1)\cdots(n-k+1)}{k!} 2^{-\binom{k}{2}} < \frac{n^k 2^{-\binom{k}{2}}}{k!}. \end{aligned}$$

# Theorem 12.12 (continued 2)

**Theorem 12.12.** For all positive integers k,  $r(k, k) \ge 2^{k/2}$ .

**Proof (continued).** Suppose  $n < 2^{k/2}$ . Then, since  $k \ge 3$ ,

$$\frac{|\mathcal{G}_n^k|}{|\mathcal{G}_n|} < \frac{n^k 2^{-\binom{k}{2}}}{k!} < \frac{2^{k^2/2} 2^{-\binom{k}{2}}}{k!} = \frac{2^{k/2}}{k!} < \frac{1}{2}.$$

That is, if  $n < 2^{k/2}$  then strictly fewer than half of the graphs in  $\mathcal{G}_n$  contain a stable set of k vertices. By considering complements, we similarly have that strictly fewer than half of the graphs in  $\mathcal{G}_n$  contain a stable set of kvertices. Therefore some graph in  $\mathcal{G}_n$  contains neither a clique of k vertices nor a stable set of k vertices. That is, if  $n < 2^{k/2}$  then there aren't necessarily enough vertices in a graph on n vertices to guarantee that the graph either contains a clique on k vertices or a stable set on k vertices. Hence, n < r(k, k) and so we must have  $r(k, k) \ge 2^{k/2}$  as claimed.

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# Theorem 12.15. Schur's Theorem

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Let  $\{A_1, A_2, \ldots, A_n\}$  be a partition of the set of integers  $\{1, 2, \ldots, r_n\}$  into n subsets. Then some  $A_i$  contains three integers x, y, and z satisfying the equation z.

**Proof.** Consider the complete graph whose vertex set is  $\{1, 2, ..., r_n\}$ . Color the edges of this graph with colors 1, 2, ..., n by the rule that the edge uv is assigned color i if  $|u - v| \in A_i$ . By the definition of this general Ramsey number  $r_n = r(t_1, t_2, ..., t_n) = r(3, 3, ..., 3)$  we know that some  $A_j$  contains a  $K_3$ ; that is, there are three vertices a, b, c such that edges ab, bc, and ac all have the same color j. Suppose, without loss of generality, that a > b > c. Let x = a - b, y = b - c, and z = a - c. Then, since ab, bc, ac are color j, then  $x, y, z \in A_j$ . Also, x + y = (a - b) + (b - c) = a - c = z, as claimed.

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