# Graph Theory

#### **Chapter 13. The Probabilistic Method** 13.5. The Local Lemma—Proofs of Theorems



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**Theorem 13.12.** THE LOCAL LEMMA.

Let  $A_i$ , where  $i \in N$ , be events in a finite probability space  $(\Omega, P)$  and let  $N_i \subseteq N$  where  $i \in N$ . Suppose that, for all  $i \in N$ ,

(i)  $A_i$  is independent of the set of events  $\{A_j \mid J \in N_i\}$ ,

(ii) for each  $i \in N$ , there is a constant  $p_i$  where  $0 < p_i < 1$ , and for each  $i \in N$  we have  $P(A_i) = p_i \prod_{j \in N_i} (1 - p_j)$ .

Set  $B_i = \overline{A}_i$  where  $i \in N$ . Then, for any two disjoint subsets  $R, S \subseteq N$ ,

$$P(B_R \cap B_S) \ge P(B_R) \prod_{i \in S} (1 - p_i). \tag{13.15}$$

In particular, when  $R = \emptyset$  and S = N,

$$P\left(\cap_{i\in N}\overline{A}_i\right) \geq \prod_{i\in N} (1-p_i) > 0.$$
 (13.16)

# Theorem 13.12 (continued 1)

**Proof.** If  $S = \emptyset$  then  $B_S = \bigcap_{i \in S} B_i = \bigcap_{i \in S} \overline{A}_i = \Omega$  (we could take this as the intersection of no sets) and  $\prod_{i \in S} (1 - p_i) = 1$  (similarly, this could be taken as the definition of a product of no numbers), so

$$P(B_R \cap B_S) = P(B_R \cap \Omega) = P(B_R) = P(B_R)(1) \ge P(B_R) \prod_{i \in S} (1 - p_i),$$

and equation (13.15) holds when  $S = \emptyset$ .

# Theorem 13.12 (continued 1)

**Proof.** If  $S = \emptyset$  then  $B_S = \bigcap_{i \in S} B_i = \bigcap_{i \in S} \overline{A}_i = \Omega$  (we could take this as the intersection of no sets) and  $\prod_{i \in S} (1 - p_i) = 1$  (similarly, this could be taken as the definition of a product of no numbers), so

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#### and equation (13.15) holds when $S = \emptyset$ .

If |S| = 1 and  $S = \{i\}$ , then  $B_S = B_i$  and  $\prod_{j \in S} (1 - p_j) = 1 - p_i$ . Define  $R_1 = R \setminus N_i$  and  $S_1 = R \cap N_i$  (so that  $R = R_1 \cup S_1$ ). Then

 $\begin{array}{ll} P(A_i \cap B_R) &\leq & P(A_i \cap B_{R_1}) \text{ since } R_1 \subseteq R \text{ and so} \\ & & B_R = \cap_{i \in R} \overline{A}_i \subseteq \cap_{i \in R_1} \overline{A}_i = B_{R_1} \text{ and } A_i \cap B_R \subseteq A_i \cap B_{R_1} \\ & = & P(A_i)P(B_{R_1}) \text{ since } A_i \text{ is independent of} \\ & & \{A_j \mid j \notin N_i\} \supseteq R_1 \text{ by hypothesis (i).} \end{array}$ 

# Theorem 13.12 (continued 1)

**Proof.** If  $S = \emptyset$  then  $B_S = \bigcap_{i \in S} B_i = \bigcap_{i \in S} \overline{A}_i = \Omega$  (we could take this as the intersection of no sets) and  $\prod_{i \in S} (1 - p_i) = 1$  (similarly, this could be taken as the definition of a product of no numbers), so

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If |S| = 1 and  $S = \{i\}$ , then  $B_S = B_i$  and  $\prod_{j \in S} (1 - p_j) = 1 - p_i$ . Define  $R_1 = R \setminus N_i$  and  $S_1 = R \cap N_i$  (so that  $R = R_1 \cup S_1$ ). Then

 $\begin{array}{ll} P(A_i \cap B_R) &\leq & P(A_i \cap B_{R_1}) \text{ since } R_1 \subseteq R \text{ and so} \\ & & B_R = \cap_{i \in R} \overline{A}_i \subseteq \cap_{i \in R_1} \overline{A}_i = B_{R_1} \text{ and } A_i \cap B_R \subseteq A_i \cap B_{R_1} \\ &= & P(A_i)P(B_{R_1}) \text{ since } A_i \text{ is independent of} \\ & & \{A_j \mid j \notin N_i\} \supseteq R_1 \text{ by hypothesis (i).} \end{array}$ 

# Theorem 13.12 (continued 2)

**Proof (continued).** Since  $S_1 \subseteq N_i \subseteq N$ , then by hypothesis (ii)

$$\mathsf{P}(\mathsf{A}_i) \leq \mathsf{p}_i \prod_{j \in \mathsf{N}_i} (1-\mathsf{p}_j) \leq \mathsf{p}_i \prod_{j \in \mathsf{S}_1} (1-\mathsf{p}_j) \qquad (**)$$

for some  $p_i$  with  $0 < p_i < 1$  and some  $p_j$  with  $0 < p_j < 1$ . In Exercise 13.5.A(i) it is to be shown by induction on  $|S_1|$  that

$$P(B_{R_1})\prod_{j\in S_1}(1-p_j) \le P(B_{R_1}\cap B_{S_1}).$$
 (†)

Therefore

$$\begin{array}{lll} P(A_i \cap B_R) & \leq & P(A_i)P(B_{R_1}) \text{ by } (*) \\ & \leq & P(B_{R_1})p_i \prod_{j \in S_1} (1-p_j) \text{ by } (**) \\ & \leq & p_i P(B_{R_1} \cap B_{S_1}) \text{ by } (\dagger), \dots \end{array}$$

# Theorem 13.12 (continued 2)

**Proof (continued).** Since  $S_1 \subseteq N_i \subseteq N$ , then by hypothesis (ii)

$$\mathsf{P}(\mathsf{A}_i) \leq \mathsf{p}_i \prod_{j \in \mathsf{N}_i} (1-\mathsf{p}_j) \leq \mathsf{p}_i \prod_{j \in \mathcal{S}_1} (1-\mathsf{p}_j) \qquad (**)$$

for some  $p_i$  with  $0 < p_i < 1$  and some  $p_j$  with  $0 < p_j < 1$ . In Exercise 13.5.A(i) it is to be shown by induction on  $|S_1|$  that

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Therefore

$$\begin{array}{lll} P(A_i \cap B_R) &\leq & P(A_i)P(B_{R_1}) \text{ by } (*) \\ &\leq & P(B_{R_1})p_i \prod_{j \in S_1} (1-p_j) \text{ by } (**) \\ &\leq & p_i P(B_{R_1} \cap B_{S_1}) \text{ by } (\dagger), \dots \end{array}$$

Theorem 13.12 (continued 3)

#### Proof (continued). and so

$$P(B_R \cap B_S) = P(B_R \cap B_i) \text{ since } S = \{i\}$$
  
=  $P(B_R) - P(B_R \cap A_i) \text{ since}$   
 $B_R = (B_R \cap A_i) \cup (B_R \cap \overline{A}_i) = (B_R \cap A_i) \cup B_R \cap B_i)$   
 $\geq P(B_R - p_i P(B_R) \text{ by } (\dagger\dagger)$   
=  $P(B_R)(1 - p_i) = P(B_R) \prod_{i \in S} (1 - p_i),$ 

#### and equation (13.15) holds when |S| = 1.

If  $|S| \ge 2$ , then let  $R_1$  and  $S_1$  be nonempty disjoint sets which partition S so that  $S = R_1 \cup S_1$ . Then

 $P(B_R \cap B_S) = P(B_R \cap B_{R_1 \cup S_1})$ 

 $= P(B_R \cap B_{R_1} \cap B_{S_1})$ by definition of *B*'s as intersections  $= P(B_{R \cup R_1} \cap B_{S_1})$ similarly. Theorem 13.12 (continued 3)

#### Proof (continued). and so

$$P(B_R \cap B_S) = P(B_R \cap B_i) \text{ since } S = \{i\}$$
  
=  $P(B_R) - P(B_R \cap A_i) \text{ since}$   
 $B_R = (B_R \cap A_i) \cup (B_R \cap \overline{A}_i) = (B_R \cap A_i) \cup B_R \cap B_i)$   
 $\geq P(B_R - p_i P(B_R) \text{ by } (\dagger\dagger)$   
=  $P(B_R)(1 - p_i) = P(B_R) \prod_{i \in S} (1 - p_i),$ 

and equation (13.15) holds when |S| = 1.

If  $|S| \ge 2$ , then let  $R_1$  and  $S_1$  be nonempty disjoint sets which partition S so that  $S = R_1 \cup S_1$ . Then

$$P(B_R \cap B_S) = P(B_R \cap B_{R_1 \cup S_1})$$
  
=  $P(B_R \cap B_{R_1} \cap B_{S_1})$  by definition of *B*'s as intersections  
=  $P(B_{R \cup R_1} \cap B_{S_1})$  similarly.

Theorem 13.12 (continued 4)

**Proof (continued).** In Exercise 13.5.A(ii) it is to be shown by induction on  $|S_1|$  that

$$P(B_{R\cup R_{1}} \cap B_{S_{1}}) \geq P(B_{R\cup R_{1}}) \prod_{i \in S_{1}} (1-p_{i})$$
  
=  $P(B_{R} \cap B_{R_{1}}) \prod_{i \in S_{1}} (1-p)i).$  (‡)

In Exercise 13.5.A(iii) it is to be shown by induction on  $|R \cup R_1|$  that

$$P(B_R \cap B_{R_1}) \ge P(B_R) \prod_{i \in R_1} (1 - p_i). \tag{\ddagger}$$

# Theorem 13.12 (continued 5)

#### **Proof (continued).** Therefore

$$P(B_R \cap B_S) = P(B_{R \cup R_1} \cap B_{S_1})$$

$$\geq P(B_R \cap B_{R_1}) \prod_{i \in S_1} (1 - p_i) \text{ by } (\ddagger)$$

$$\geq P(B_R) \prod_{i \in R_1} (1 - p_i) \prod_{i \in S_1} (1 - p_i) \text{ by } (\ddagger\ddagger)$$

$$= P(B_R) \prod_{i \in S} (1 - p_i) \text{ since } S = R_1 \cup S_1,$$

so equation (13.15) holds when  $|S| \ge 2$  and hence holds for all S and R subsets of N, as claimed.

**Theorem 13.14.** THE LOCAL LEMMA—SYMMETRIC VERSION. let  $A_i$ , where  $i \in N$ , be events in a finite probability space  $(\Omega, P)$  having a dependency graph with maximum degree d. Suppose  $P(A_i) < 1/(e(d+1))$  for all  $i \in N$  (where "e" here is the base of the natural log function). Then  $P(\cap_{i \in N} \overline{A_i}) > 0$ .

**Proof.** Set  $p_1 = 1/(d+1) = p$  for  $i \in N$  (this value of p maximizes the function  $f(p) = p(1-p)^d$  for  $p \in (0,1)$  and will give us a "uniform bound" on  $p(A_i)$  in hypothesis (ii) of The Local Lemma). Now the sets  $N_i$  are defined from the dependency graph ( $N_i$  includes all neighbors of vertex i in the dependency graph, so we have event  $A_i$  is independent of the events  $\{A_i \mid j \notin N_i\}$ , as required by hypothesis (i) of The Local Lemma).

**Theorem 13.14.** THE LOCAL LEMMA—SYMMETRIC VERSION. let  $A_i$ , where  $i \in N$ , be events in a finite probability space  $(\Omega, P)$  having a dependency graph with maximum degree d. Suppose  $P(A_i) < 1/(e(d+1))$  for all  $i \in N$  (where "e" here is the base of the natural log function). Then  $P(\cap_{i \in N} \overline{A_i}) > 0$ .

**Proof.** Set  $p_1 = 1/(d+1) = p$  for  $i \in N$  (this value of p maximizes the function  $f(p) = p(1-p)^d$  for  $p \in (0,1)$  and will give us a "uniform bound" on  $p(A_i)$  in hypothesis (ii) of The Local Lemma). Now the sets  $N_i$  are defined from the dependency graph ( $N_i$  includes all neighbors of vertex i in the dependency graph, so we have event  $A_i$  is independent of the events  $\{A_j \mid j \notin N_i\}$ , as required by hypothesis (i) of The Local Lemma).

Theorem 13.14. The Local Lemma—Symmetric Version

# Theorem 13.14 (continued)

#### Proof (continued). Now

$$\begin{array}{ll} \mathcal{P}(A_i) &\leq \quad \displaystyle\frac{1}{e(d+1)} \text{ by hypothesis} \\ &\leq \quad \left(\displaystyle\frac{d}{d+1}\right)^d \left(\displaystyle\frac{1}{d+1}\right) \text{ since } 1 + \displaystyle\frac{1}{d} \leq e^{1/d} \text{ by Exercise 13.2.1(b)} \\ &\quad \text{with } x = 1/d, \text{ or } (1+1/d)^d \leq e \text{ or } \left(\displaystyle\frac{d+1}{d}\right)^d \leq e \\ &\quad \text{ or } \left(\displaystyle\frac{d}{d+1}\right)^d \geq \displaystyle\frac{1}{e} \\ &= \quad p \prod_{j \in N_i} (1-p)^d \text{ since } p_i = p = \displaystyle\frac{1}{d+1} \text{ for all } i \in N. \end{array}$$

So hypothesis (ii) of the Local Lemma holds. Hence, by the Local Lemma (Theorem 13.12, the "in particular" part),  $P(\bigcap_{i \in N} \overline{A}_i) > 0$ , as claimed.  $\Box$ 

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**Theorem 13.15.** Let  $H = (V, \mathcal{F})$  be a hypergraph in which each edge has at least k elements and meets at most d other edges. If  $e(d+1) \leq 2^{k-1}$  (again, "e" here is the base of the natural log function), then H is 2-colourable.

**Proof.** Consider a random 2-colouring of V, where each vertex receives one of two colours with probability 1/2. For each edge F of H, denote by  $A_F$  the event that F is monochomatic. Then events  $A_F$  and  $A_G$  are independent unless edges F and G share vertices. Since an edge of Hmeets at most d other edges, then the dependence graph for the events  $\{A_F \mid F \in \mathcal{F}\}$  has maximum degree d.

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**Theorem 13.15.** Let  $H = (V, \mathcal{F})$  be a hypergraph in which each edge has at least k elements and meets at most d other edges. If  $e(d+1) \leq 2^{k-1}$  (again, "e" here is the base of the natural log function), then H is 2-colourable.

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# **Corollary 13.16.** Let $H = (V, \mathcal{F})$ be a *k*-uniform *k*-regular hypergraph, where $k \ge 9$ . Then *H* is 2-colourable.

**Proof.** Since *H* is *k*-uniform, then each edge contains exactly *k* elements (i.e., vertices), and since *H* is *k*-regular then each vertex of *H* lies on *k* edges. So for given edge *F* of *H*, *F* contains *k* vertices and each lies on k - 1 edges (along with edge *F*), so that edge *F* meets at most d = k(k - 1) other edges.

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**Corollary 13.16.** Let  $H = (V, \mathcal{F})$  be a *k*-uniform *k*-regular hypergraph, where  $k \ge 9$ . Then *H* is 2-colourable.

**Proof.** Since *H* is *k*-uniform, then each edge contains exactly *k* elements (i.e., vertices), and since H is k-regular then each vertex of H lies on kedges. So for given edge F of H, F contains k vertices and each lies on k-1 edges (along with edge F), so that edge F meets at most d = k(k-1) other edges. If k > 9 then  $e(d+1) = e(k(k-1)+1) \le 2^{k-1}$  since (1) for k = 9 we have  $e(k(k-1)+1) = 73e \approx 198.4, 2^{k-1} = 256$ , and (2)  $f(x) = 2^{x-1} - e(x^2 - x + 1)$  has derivative  $f'(x) = (\ln 2)2^{x-1} - e(2x-1) > 0$  for x > 9, so that f is increasing for  $x \ge 9$  and hence  $f(x) \ge f(9) > 0$  for all  $x \ge 9$ . So the hypotheses of Theorem 13.15 are satisfied and hence H is 2-colourable.

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# **Theorem 13.17.** Let *D* be a strict (i.e., "simple") *k*-diregular digraph where $k \ge 8$ . Then *D* contains a directed even cycle.

**Proof.** Consider a random 2-colouring of V, where each vertex receives one of the two colours with probability 1/2. For each vertex v of D, denote by  $A_v$  the event that c(u) = c(v) for all  $u \in N^+(v)$  (that is,  $A_v$  denotes the event that all outneighbors of v are the same colour as v). So  $A_v$  is independent of all  $A_u$  such that  $(\{u\} \cup N^+(u)) \cap N^+(v) = \emptyset$  (that is, the outneighbors of v do not include any outneighbors of u nor u itself). Then  $A_v$  is dependent on some  $A_u$  when onr (or more) of the k outneighbors of v is one of the k outneighbors of u or u itself.

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**Proof.** Consider a random 2-colouring of V, where each vertex receives one of the two colours with probability 1/2. For each vertex v of D, denote by  $A_v$  the event that c(u) = c(v) for all  $u \in N^+(v)$  (that is,  $A_v$  denotes the event that all outneighbors of v are the same colour as v). So  $A_v$  is independent of all  $A_u$  such that  $(\{u\} \cup N^+(u)) \cap N^+(v) = \emptyset$  (that is, the outneighbors of v do not include any outneighbors of u nor u itself). Then  $A_{v}$  is dependent on some  $A_{u}$  when onr (or more) of the k outneighbors of v is one of the k outneighbors of u or u itself. Since D is k-diregular, then each of the k-outneighbors of v can be an outneighbor of k-1 other vertices (along with being an outneighbor of v itself) and is a vertex associated with some  $A_{\mu}$ , so that each outneighbor of v is associated with up to k other events dependent on  $A_{\nu}$ . So in the dependency graph of the event  $\{A_v \mid v \in V\}$ , each vertex is of degree at most  $d = k^2$ .

**Theorem 13.17.** Let *D* be a strict (i.e., "simple") *k*-diregular digraph where  $k \ge 8$ . Then *D* contains a directed even cycle.

**Proof.** Consider a random 2-colouring of V, where each vertex receives one of the two colours with probability 1/2. For each vertex v of D, denote by  $A_v$  the event that c(u) = c(v) for all  $u \in N^+(v)$  (that is,  $A_v$  denotes the event that all outneighbors of v are the same colour as v). So  $A_v$  is independent of all  $A_u$  such that  $(\{u\} \cup N^+(u)) \cap N^+(v) = \emptyset$  (that is, the outneighbors of v do not include any outneighbors of u nor u itself). Then  $A_v$  is dependent on some  $A_u$  when onr (or more) of the k outneighbors of v is one of the k outneighbors of u or u itself. Since D is k-diregular, then each of the k-outneighbors of v can be an outneighbor of k-1 other vertices (along with being an outneighbor of v itself) and is a vertex associated with some  $A_{\mu}$ , so that each outneighbor of v is associated with up to k other events dependent on  $A_{v}$ . So in the dependency graph of the event  $\{A_v \mid v \in V\}$ , each vertex is of degree at most  $d = k^2$ .

# Theorem 13.17 (continued 1)

**Theorem 13.17.** Let *D* be a strict (i.e., "simple") *k*-diregular digraph where  $k \ge 8$ . Then *D* contains a directed even cycle.

**Proof (continued).** Next, notice that  $\overline{A}_v$  is the event that v has an outneighbor of a different colour. Since v has k outneighbors then  $P(A_v) = 1/2^k$ . With  $d = k^2$  we have  $P(A_i) = \frac{1}{2^k} \le \frac{1}{e(d+1)} = \frac{1}{e(k^2+1)}$  for  $k \ge 8$ , since (1) for k = 8 we have  $\frac{1}{2^k} = \frac{1}{2^8} = \frac{1}{256} \approx 0.0039$ ,  $\frac{1}{e(k^2+1)} = \frac{1}{65e} \approx 0.0057$ , and (2)  $f(x) = \frac{1}{e(x^2+1)} - 2^{-x}$  has derivative  $f'(x) = \frac{-2x}{e(x^2+1)^2} + (\ln 2)2^{-x} > 0$ for  $k \ge 8$  so that f is increasing for  $x \ge 8$  and hence  $f(x) \ge f(8) > 0$  for all x > 8. So the hypotheses of Theorem 3.14 are satisfied and hence  $P(\bigcap_{v \in V} \overline{A}_v) > 0$ . That is, there is a 2-colouring of V such that every  $v \in V$  has an outneighbor of the opposite colour.

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# Theorem 13.17 (continued 2)

**Theorem 13.17.** Let *D* be a strict (i.e., "simple") *k*-diregular digraph where  $k \ge 8$ . Then *D* contains a directed even cycle.

**Proof (continued).** With respect to this colouring, let uPv be a maximal (length) properly 2-coloured directed path in D and let w be an outneighbor of v of the opposite colour of v. Since uPv is maximal, then w must be some vertex of uPv (or else uPvw would be a longer properly 2-coloured path in D). So take the segment of P from w to v (denoted wPv) and then add the arc from v to w to produce a cycle in D. This cycle is properly coloured and so is an even length cycle, as claimed.



#### Lemma 13.18

**Lemma 13.18.** Let G = (V, E) be a simple graph and let  $\{V_1, V_2, \ldots, V_k\}$  be a partition of V into k sets, each of cardinality at least  $2e\Delta$  (again, "e" here is the base of the natural log function). Then there is a stable set S in G such that  $|S \cap V_i| = 1$  for  $1 \le i \le k$ .

**Proof.** By deleting vertices from *G* if necessary, we may assume that  $|V_i| = t = \lceil 2e\Delta \rceil$  for  $1 \le i \le k$  (we'll show the existence of a stable set *S* under these conditions, then the deleted vertices and relevant edges can be added back in to *G* and this won't have an effect on the stability of set *S* nor on the intersection property of set *S*). We select one vertex  $v_i$  at random from each set  $V_i$  for  $1 \le i \le k$ , and set  $S = \{v_1, v_2, \ldots, v_n\}$ .

#### Lemma 13.18

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For an edge f of G, let  $A_f$  denote the event that both ends of f belong to S. Since  $|V_i| = t$  for each i then  $P(A_f) = 1/t^2$  for all  $f \in E$ . In Exercise 13.5.A is to be shown that  $A_f$  is dependent only on those events  $A_g$  such that an end of g lies in the same set  $V_i$  as an end of f.

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**Proof (continued).** For a given  $A_f$  where f has its ends in  $V_i$  and  $V_j$  (where possibly i = j), there are at most  $t\Delta$  edges with an end in  $V_i$  and at most  $t\Delta$  edges with an end in  $V_j$ , so that there are at most  $2t\Delta - 1$  other events  $A_g$  which are dependent on  $A_f$ . We set  $d = 2t\Delta - 1$  so that in the dependency graph for events  $\{A_f \mid f \in E\}$  has maximum degree at most  $d = 2t\Delta - 1$ . Also,

$$P(A_f) = \frac{1}{t^2} = \frac{1}{t\lceil 2e\Delta\rceil} \le \frac{1}{t(2e\Delta)}$$
$$= \frac{1}{e(2t\Delta)} = \frac{1}{e((2t\Delta - 1) + 1)} = \frac{1}{e(d+1)}.$$

# Lemma 13.18 (continued 1)

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# Lemma 13.18 (continued 2)

**Lemma 13.18.** Let G = (V, E) be a simple graph and let  $\{V_1, V_2, \ldots, V_k\}$  be a partition of V into k sets, each of cardinality at least  $2e\Delta$  (again, "e" here is the base of the natural log function). Then there is a stable set S in G such that  $|S \cap V_i| = 1$  for  $1 \le i \le k$ .

**Proof (continued).** So the hypotheses of Theorem 13.14 hold and hence  $P(\bigcap_{f \in E} \overline{A}_f) > 0$ . So by Note 13.5.A, there exists a set *S* which (by construction) intersects each  $V_i$  in one point and each edge *f* of *G* has its ends in different sets  $V_i$  and  $V_j$  (since  $\overline{A}_f$  holds; i.e., *f* does not have both ends in the same  $V_i$ ). Since *S* contains exactly one point from each  $V_i$ , then every edge incident to a vertex in *S* has its other end outside of *S*. That is, *S* is a stable set, as claimed.

**Theorem 13.19.** Let G = (V, E) be a simple 2*r*-regular graph with girth at least 2e(4r - 2) (again, "e" here is the base of the natural log function). Then la(G) = r + 1.

**Proof.** We saw in Note 13.5.B that  $la(G) \le r + 1$ . We now borrow a result from Section 16.4 ("Perfect Matchings and Factors"): "Every 2*r*-regular graph admits a decomposition into 2-factors" (this is Exercise 16.4.16).

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Consider such a decomposition  $\{F_1, F_2, \ldots, F_r\}$  of G and let  $C_i$ , for  $1 \le i \le k$ , be the constituent cycles of these 2-factors (since a 2-factor is a 2-regular graph, then it is a vertex disjoint union of cycles). Define the *edge* sets  $V_i = E(C_i)$  for  $1 \le i \le k$ . The line graph H of G is (4r - 2)-regular (since each edge of G is adjacent to 2r - 1 other edges at each of its ends).

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**Proof (continued).** Because G has girth at least 2e(4r-2) by hypothesis, then the edge sets  $\{V_1, V_2, \ldots, V_k\}$  partition the vertex set V(H) into k sets, each of cardinality at least 2e(4r-2) (since each is the set of edges of a cycle in G). So the hypotheses of Lemma 13.18 are satisfied by the line graph H, so H has a stable set S meeting each set  $V_i$ in one vertex. Define the subgraphs  $L_i = F_i \setminus S$  for  $1 \le i \le r$ . Since  $F_i$  is a collection of vertex disjoint cycles in G and S includes an edge of each of these cycles, then each  $L_i$  is a linear forest on G. Also, if we set  $L_0$ equal to the subgraph of G which has edge set S and vertex set of all ends of edges in S (so that  $L_0$  is a linear forest where each tree has one edge because S is a stable set). Then  $\{L_0, L_1, \ldots, L_r\}$  is a decomposition of G into r + 1 linear forests. So la(G) = r + 1.

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