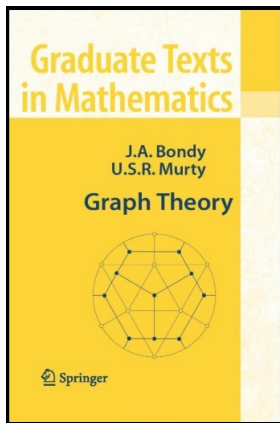


# Graph Theory

## Chapter 14. Vertex Colourings

### 14.2. Critical Graphs—Proofs of Theorems



# Table of contents

1 Theorem 14.7

2 Theorem 14.8

3 Theorem 14.10

# Theorem 14.7

**Theorem 14.7.** If  $G$  is  $k$ -critical then  $\delta \geq k - 1$ .

**Proof.** Let  $G$  be  $k$ -critical. ASSUME  $\delta < k - 1$ . Let  $v$  be a vertex of degree  $\delta$  in  $G$ . Since  $G$  is  $k$ -critical, then  $G - v$  is  $(k - 1)$ -colourable (since we are removing only one vertex then the number of colours required to properly colour  $G - v$  is at most one less; since  $G$  is  $k$ -critical, it is exactly one less). Let  $\{V_1, V_2, \dots, V_{k-1}\}$  be the colour classes of a  $(k - 1)$ -colouring of  $G - v$ .

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**Theorem 14.8.** No critical graph has a clique cut.

**Proof.** Let  $G$  be a  $k$ -critical graph. ASSUME that  $G$  has a clique cut  $S$ . Denote the  $S$ -components of  $G$  by  $G_1, G_2, \dots, G_t$ . Since  $G$  is  $k$ -critical, each  $G_i$  is  $(k - 1)$ -colourable (though maybe not  $(k - 1)$ -chromatic). Since  $S$  is assumed to be a clique, then the vertices of  $S$  receive distinct colours in any  $(k - 1)$ -colouring of  $G_i$ .

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**Theorem 14.10.** Let  $G$  be a  $k$ -critical graph with a 2-vertex cut set  $\{u, v\}$ , and let  $e$  be a new edge joining  $u$  and  $v$ . Then

- (1)  $G = G_1 \cup G_2$ , where  $G_i$  is a  $\{u, v\}$ -component of  $G$  of Type  $i$  for  $i \in \{1, 2\}$ ,
- (2) both  $H_1 = G_1 + e$  and  $H_2 = G_2/\{u, v\}$  are  $k$  critical.

**Proof.** For (1), because  $G$  is  $k$ -critical then each  $\{u, v\}$ -component of  $G$  is  $(k - 1)$ -colourable. There cannot be  $(k - 1)$ -colourings of the  $\{u, v\}$ -components all of which agree on  $\{u, v\}$ , since this would imply a  $(k - 1)$ -colouring of  $G$ . Therefore there are two  $\{u, v\}$ -components  $G_1$  and  $G_2$  such that no  $(k - 1)$ -colouring of  $G_1$  agrees with any  $(k - 1)$ -colouring of  $G_2$ .

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**Proof (continued).** Since  $G$  is  $k$ -critical then we must have  $G = G_1 \cup G_2$ , as claimed. (Notice that  $G = G_1 \cup G_2$  implies that  $G$  does not have an edge joining  $u$  and  $v$ .)

For (2), because  $G_1$  is of Type 1,  $H_1 = G_1 + e$  is  $k$ -chromatic (since a  $(k - 1)$ -colouring of  $G_1$  only exists when  $u$  and  $v$  are the same colour). Let  $f$  be any edge of  $H_1$  (we show that  $H_1 \setminus f$  is  $(k - 1)$ -colourable, and hence  $H_1$  is  $k$ -critical). If  $f = e$  then  $H_1 \setminus f = H_1 \setminus e = G_1$  and so  $H_1 \setminus f$  is  $(k - 1)$ -colourable. Let  $f$  be any edge of  $H_1$  other than  $e$  (so  $f$  is an edge of  $G_1$ ). Any  $(k - 1)$ -colouring of  $G \setminus f$  yields a  $(k - 1)$ -colouring of  $G_2$  and so  $u$  and  $v$  must receive different colours (since  $G_2$  is Type 2 by (1)).

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