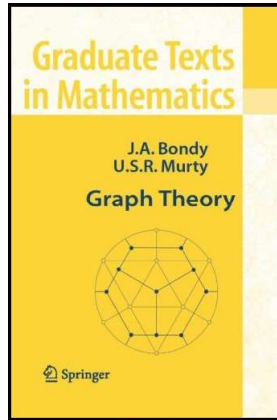


# Graph Theory

## Chapter 14. Vertex Colourings

### 14.4. Perfect Graphs—Proofs of Theorems



## Proposition 14.15

**Proposition 14.15.** Let  $S$  be a stable set in a minimally imperfect graph  $G$ . Then  $\omega(G - S) = \omega(G)$ .

**Proof.** This is easy with the use of Exercise 14.4.5, which states that for  $S$  a stable set in a minimally imperfect graph  $G$  we have

$$\omega(G - S) \leq \omega(G) \leq \chi(G) - 1 \leq \chi(G - S) = \omega(G - S).$$

Since the left and right terms are the same, then the inequalities reduce to equalities.  $\square$

## Lemma 14.16

**Lemma 14.16.** Let  $G$  be a minimally imperfect graph with stability number  $\alpha$  and clique number  $\omega$ . Then  $G$  contains  $\alpha\omega + 1$  stable sets  $S_0, S_1, \dots, S_{\alpha\omega}$  and  $\alpha\omega + 1$  cliques  $C_0, C_1, \dots, C_{\alpha\omega}$  such that:

- each vertex of  $G$  belongs to precisely  $\alpha$  of the stable sets  $S_i$ ,
- each clique  $C_j$  has  $\omega$  vertices,
- $C_j \cap S_i = \emptyset$  for  $0 \leq i \leq \alpha\omega$ , and
- $|C_j \cap S_i| = 1$  for  $0 \leq i < j \leq \alpha\omega$ .

**Proof.** Let  $S_0$  be a stable set of  $\alpha$  vertices of  $G$ , and let  $v \in S_0$  (of course  $|S_0| \geq 1$ ). The graph  $G - v$  is perfect because  $G$  is minimally imperfect (note that  $G - v$  is the subgraph of  $G$  induced by  $V(G) \setminus \{v\}$ ). So by the definition of “perfect graph” and Note 14.1.B,  $\chi(G - v) = \omega(G - v) \leq \omega(G)$ . Since the colour classes of a proper colouring are each stable sets, we then have for any  $v \in S_0$  that the set  $V \setminus \{v\}$  can be partitioned into a family  $\mathcal{S}_v$  of  $\omega$  stable sets (the inequality tells us that there are at most  $\omega$  stable sets, but the stable sets can be further subdivided to get a total of  $\omega$  stable sets).

## Lemma 14.16 (continued 1)

**Proof (continued).** Since  $|S_0| = \alpha$  then there are  $\alpha$  such  $v$  in  $S_0$ , since each family  $\mathcal{S}_v$  contains  $\omega$  stable sets, then there are  $\alpha\omega$  stable sets in the totality of the families  $\mathcal{S}_1, \mathcal{S}_2, \dots, \mathcal{S}_\alpha$ . Denote these stable sets as  $S_1, S_2, \dots, S_{\alpha\omega}$ . We'll see below that each  $S_j$  is nonempty when we establish the fourth claim. For each  $v$  a vertex of  $G$  where  $v \in S_0$ , we have that  $v \in S_i$  for some unique  $S_i$  in family  $\mathcal{S}_w$  where  $w \neq v$ . So this  $v$  is in  $S_0$  and all but one of  $S_1, S_2, \dots, S_\alpha$ ; that is,  $v$  is in  $\alpha$  of the stable sets  $S_j$ . For each  $v$  a vertex of  $G$  where  $v \notin S_0$ , we have that  $v$  is in some unique  $S_i$  in family  $\mathcal{S}_j$  for each  $1 \leq j \leq \alpha$ ; that is,  $v$  is in  $\alpha$  of the stable sets  $S_j$ . Therefore the first claim holds.

By Proposition 14.15, the clique numbers satisfy  $\omega(G - S_i) = \omega(G)$  for each  $0 \leq i \leq \alpha\omega$ . Therefore there is a maximum clique  $C_i$  of  $G - S_i$  that is also a maximum clique of  $G$ . Each  $C_i$  has  $\omega = \omega(G)$  vertices, so the second claim holds. So there is a maximum clique  $C_j$  of  $G$  that is disjoint from  $S_i$  and the third claim holds.

## Lemma 14.16 (continued 2)

**Lemma 14.16.** Let  $G$  be a minimally imperfect graph with stability number  $\alpha$  and clique number  $\omega$ . Then  $G$  contains  $\alpha\omega + 1$  stable sets  $S_0, S_1, \dots, S_{\alpha\omega}$  and  $\alpha\omega + 1$  cliques  $C_0, C_1, \dots, C_{\alpha\omega}$  such that:

- each vertex of  $G$  belongs to precisely  $\alpha$  of the stable sets  $S_i$ ,
- each clique  $C_i$  has  $\omega$  vertices,
- $C_i \cap S_i = \emptyset$  for  $0 \leq i \leq \alpha\omega$ , and
- $|C_i \cap S_j| = 1$  for  $0 \leq i < j \leq \alpha\omega$ .

**Proof (continued).** Because each of the  $\omega$  vertices of  $C_i$  lie in  $\alpha$  of the stable sets (by the first claim), there are  $\alpha\omega + 1$  stable sets (but  $C_i \cap S_i = \emptyset$  for  $0 \leq i \leq \alpha\omega$  by the third claim), and no two vertices of clique  $C_i$  can belong to a common stable set (since vertices of a clique are adjacent and vertices of a stable set are not adjacent), then each  $C_i$  shares exactly one point with each  $S_j$  where  $i \neq j$ . That is,  $|C_i \cap S_j| = 1$  for  $0 \leq i < j \leq \alpha\omega$ , and the fourth claim holds.  $\square$

## Theorem 14.14

**Theorem 14.14.** A graph  $G$  is perfect if and only if every induced subgraph  $H$  of  $G$  satisfies the inequality  $v(H) \leq \alpha(H)\omega(H)$ .

**Proof.** Suppose that  $G$  is a perfect graph so that  $\chi(H) = \omega(H)$  for all induced subgraphs  $H$  of  $G$ . Let  $H$  be any (fixed) induced subgraph of  $G$ . Then  $H$  is  $\omega(H)$ -colourable (since  $\chi(H) = \omega(H)$ ). Since a colour class is an independent set and the largest independent set is of size  $\alpha$ , then the number of vertices satisfies  $v(H) \leq \alpha(H)\chi(H) = \alpha(H)\omega(H)$ . So if  $G$  is perfect, then for every induced subgraph  $H$  of  $G$  we have  $v(H) \leq \alpha(H)\omega(H)$ .

We now prove the converse (actually, the contrapositive of the converse). Suppose  $G$  is not perfect. We want to show that  $v(G) \geq \alpha(G)\omega(G) + 1$ . If we can show this for minimally imperfect graphs, then it will hold for all imperfect graphs. So without loss of generality, suppose  $G$  is minimally imperfect.

## Theorem 14.14 (continued 1)

**Proof (continued).** Consider the stable sets  $S_i$  for  $0 \leq i \leq \alpha\omega$  and the cliques  $C_i$  for  $0 \leq i \leq \alpha\omega$  as given for a minimally imperfect graph in Lemma 14.16. Let  $\mathbf{S}$  and  $\mathbf{C}$  be the  $n \times (\alpha\omega + 1)$  incidence matrices of the families. Now the transpose of  $\mathbf{S}$ ,  $\mathbf{S}^t$ , is  $(\alpha\omega + 1) \times n$  and so the product  $\mathbf{S}^t\mathbf{C}$  exists and is  $(\alpha\omega + 1) \times (\alpha\omega + 1)$ . Now the  $i$ th row of  $\mathbf{S}^t$  consists of 0's and 1's, where the  $k$ th entry is 1 if  $k \in S_i$  and 0 if  $k \notin S_i$ . The  $j$ th column of  $\mathbf{C}$  consists of 0's and 1's where the  $k$ th entry is 1 if  $k \in C_j$  and 0 if  $k \notin C_j$ . The  $(i, j)$ -entry of  $\mathbf{S}^t\mathbf{C}$  is the inner product (or "dot product") of the  $i$ th row of  $\mathbf{S}^t$  and the  $j$ th column of  $\mathbf{C}$ . By the third claim of Lemma 14.16 we have  $C_i \cap S_i = \emptyset$  for  $0 \leq i \leq \alpha\omega$ , so in the  $i$ th row of  $\mathbf{S}^t$  and the  $i$ th column of  $\mathbf{C}$ , if the  $k$ th entry of one of these is 1 then the  $k$ th entry of the other is 0. So the inner product is 0 and the  $(i, i)$ -entry of  $\mathbf{S}^t\mathbf{C}$  is 0. By the fourth claim of Lemma 14.16 we have  $|C_i \cap S_j| = 1$  for  $0 \leq i < j \leq \alpha\omega$ , so in the  $i$ th row of  $\mathbf{S}^t$  and the  $j$ th column of  $\mathbf{C}$  the  $k$  entry of both is 1 only once. So the inner product is 1 and the  $(i, j)$ -entry of  $\mathbf{S}^t\mathbf{C}$  is 1 for  $i \neq j$ .

## Theorem 14.14 (continued 2)

**Proof (continued).** Therefore,  $\mathbf{S}^t\mathbf{C} = \mathbf{J} - \mathbf{I}$  where  $\mathbf{J}$  is the  $(\alpha\omega + 1) \times (\alpha\omega + 1)$  matrix of all 1's and  $\mathbf{I}$  is the  $(\alpha\omega + 1) \times (\alpha\omega + 1)$  identity matrix. Now  $\mathbf{J} - \mathbf{I}$  is nonsingular (i.e., invertible) because the inverse is  $(1/(\alpha\omega))\mathbf{J} - \mathbf{I}$  because

$$\begin{aligned} (\mathbf{J} - \mathbf{I}) \left( \frac{1}{\alpha\omega} \mathbf{J} - \mathbf{I} \right) &= \frac{1}{\alpha\omega} \mathbf{J}^2 - \mathbf{J} - \frac{1}{\alpha\omega} \mathbf{J} + \mathbf{I}^2 = \frac{1}{\alpha\omega} \mathbf{J}^2 - \left( 1 + \frac{1}{\alpha\omega} \right) \mathbf{J} + \mathbf{I} \\ &= \frac{1}{\alpha\omega} ((\alpha\omega + 1)\mathbf{J}) - \left( \frac{\alpha\omega + 1}{\alpha\omega} \right) \mathbf{J} + \mathbf{I} = \frac{\alpha\omega + 1}{\alpha\omega} \mathbf{J} - \frac{\alpha\omega + 1}{\alpha\omega} \mathbf{J} + \mathbf{I} = \mathbf{I}. \end{aligned}$$

This shows that  $(1/(\alpha\omega))\mathbf{J} - \mathbf{I}$  is a right inverse of  $\mathbf{J} - \mathbf{I}$ , but since  $\mathbf{I}$  and  $\mathbf{J}$  commute, this is sufficient to establish that we have a two-sided inverse (see also "Theorem 1.11. A Commutative Property" in my Linear Algebra [MATH 2010] online notes on [Section 1.5. Inverses of Matrices, and Linear Systems](#)). Since  $\mathbf{J} - \mathbf{I}$  is invertible, then it must be a full-rank  $\alpha\omega + 1$  (see "Theorem 2.6. An Invertible Criteria" in my Linear Algebra online notes on [Section 2.2. The Rank of a Matrix](#)).

## Theorem 14.14 (continued 3)

**Theorem 14.14.** A graph  $G$  is perfect if and only if every induced subgraph  $H$  of  $G$  satisfies the inequality  $v(H) \leq \alpha(H)\omega(H)$ .

**Proof (continued).** Recall that the rank of a matrix in the (common) dimension of the row space and the column space. In general,  $\text{rank}(AB) \leq \min\{\text{rank}(A), \text{rank}(B)\}$  (see Exercises 2.2.18 and 2.2.20 in J. Fraleigh and R. Beauregard's *Linear Algebra*, 3rd Edition, Addison-Wesley (1994), or see Theorem 3.3.5 in my online notes for Theory of Matrices [MATH 5090] on [Section 3.3. Matrix Rank and the Inverse of a Full Rank Matrix](#)) so

$$\alpha\omega + 1 = \text{rank}(\mathbf{J} - \mathbf{I}) = \text{rank}(\mathbf{S}^t\mathbf{C}) \leq \min\{\text{rank}(\mathbf{S}^t), \text{rank}(\mathbf{C})\}.$$

Now the row rank equals the column rank for any matrix (see "Theorem 2.4. Row Rank Equals Column Rank" in my Linear Algebra notes on [Section 2.2. The Rank of a Matrix](#)), so we have  $\text{rank}(\mathbf{S}) = \alpha\omega + 1$  and  $\text{rank}(\mathbf{C}) = \alpha\omega + 1$  since the rank can't be greater than the number of rows or the number of columns.

## Theorem 14.14 (continued 4)

**Theorem 14.14.** A graph  $G$  is perfect if and only if every induced subgraph  $H$  of  $G$  satisfies the inequality  $v(H) \leq \alpha(H)\omega(H)$ .

**Proof (continued).**

Both  $\mathbf{S}$  and  $\mathbf{C}$  have  $n$  row, so the number of rows must be at least as big as the rank, hence  $n \geq \alpha\omega + 1$ . That is, for minimally imperfect graph  $G$ ,  $v(G) \geq \alpha(G)\omega(G) + 1$ . As explained above, the general claim now holds. □