Graph Theory

Chapter 14. Vertex Colourings 14.7. The Chromatic Polynomial—Proofs of Theorems



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$$C(G,x) = C(G \setminus e, x) - C(G/e, x).$$
(14.6)

The polynomial P(G, x) is of degree *n*, with integer coefficients which alternate in sign, leading term x^n , and constant term 0.

Proof. We give an inductive proof on m, the number of edges of G. If m = 0 then $C(G, k) = k^n$, as described in Note 1.7.A, and the polynomial $P(G, x) = x^n$ satisfies the conditions of the theorem. This establishes the base case.

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For the induction hypothesis, suppose the theorem holds for all graphs with fewer than m edges, where $m \ge 1$. Let G be a loopless graph with m edges.

Theorem 14.26 (continued 1)

Proof (continued). If G is not simple, define P(G, x) = P(H, x) where H satisfies the first claim, that is there is a polynomial P(H, x) such that P(H, x) = C(H, k) for all nonnegative integers k, the (loopless) graph G also satisfies this condition since any proper colouring of H is also a proper colouring of G (and conversely); we do not apply the operation of edge contraction here to a graph with multiple edges (this could introduce **loops).** If G is simple, let e be an edge of G. Both $G \setminus e$ and G/e have m-1 edges (in $G \setminus e$ edge e is deleted, in G/e edge e is contracted) and are loopless. By the induction hypothesis, there exists polynomials $P(G \setminus e, x)$ and P(G/e, x) such that for all nonnegative integers k, we

$$P(G \setminus e, k) = C(G \setminus e, k) \text{ and } P(G/e, k) = c(G/e, k).$$
(14.7)

The induction hypothesis also gives the from of these polynomials.

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$$P(G \setminus e, k) = C(G \setminus e, k)$$
 and $P(G/e, k) = c(G/e, k).$ (14.7)

The induction hypothesis also gives the from of these polynomials.

Theorem 14.26 (continued 2)

Proof (continued). There are nonnegative integer coefficients $a_1, a_2, \ldots, a_{n-1}$ and $b_1, b_2, \ldots, b_{n-1}$ with

$$P(G \setminus e, x) = \sum_{i=1}^{n-1} (-1)^{n-i} a_i x^i + x^n \& P(G/e, x) = \sum_{i=1}^{n-1} (-1)^{n-i-1} b_i x^i$$
(14.8)

(notice that $G \setminus e$ has n vertices, but G/e has n-1 vertices; by the induction hypothesis we have $b_{n-1} = 1$). Define $P(G, x) = P(G \setminus e, x) - P(G/e, x)$ so that P(G, x) satisfies the recursion relation (14.6). We now have

$$P(G,k) = P(G \setminus e, k) - P(G/e, k) \text{ by (14.6) with } x = k$$

= $C(G \setminus e, k) - C(G/e, k) \text{ by (14.7)}$
= $C(G, k) \text{ by (14.5).}$

So P(G, k) = C(G, k) for nonnegative integer k, as claimed.

Theorem 14.26 (continued 3)

Proof (continued). Now the form of P(G, x) is:

$$P(G, x) = P(G \setminus e, x) - P(G/e, x) \text{ by our definition of}$$

$$P(G, x) \text{ in this proof}$$

$$= \sum_{i=1}^{n-1} (-1)^{n-i} a_i x^i + x^n - \sum_{i=1}^{n-1} (-1)^{n-i-1} b_i x^i \text{ by (14.8)}$$

$$= \sum_{i=1}^{n-1} (-1)^{n-i} (a_i + b_i) x^i + x^n,$$

so that P(G, x) is of degree n, with integer coefficients which alternate in sign (notice that $a_i + b_i$ is a nonnegative integer for each i), leading term x^n , and constant term 0. This establishes the induction step. Hence, by mathematical induction, the result holds for all m and hence for all loopless graphs, or loopless simple graphs, as the hypothesis require.

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