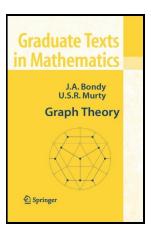
## Graph Theory

#### **Chapter 15. Colourings of Maps** 15.4. Hadwiger's Conjecture—Proofs of Theorems







#### **Theorem 15.12.** Every 4-chromatic graph contains a K<sub>4</sub>-subdivision.

**Proof.** Let *G* be a 4-chromatic graph. Then *G* contains a 4-critical subgraph *F*. By Theorem 14.7,  $\delta(F) \ge 3$ . By Exercise 10.1.5, *F* contains a subdivision of  $K_4$ . This subdivision of  $K_4$  is also a subgraph of *G*, as claimed.

Graph Theory

**Theorem 15.12.** Every 4-chromatic graph contains a  $K_4$ -subdivision.

**Proof.** Let *G* be a 4-chromatic graph. Then *G* contains a 4-critical subgraph *F*. By Theorem 14.7,  $\delta(F) \ge 3$ . By Exercise 10.1.5, *F* contains a subdivision of  $K_4$ . This subdivision of  $K_4$  is also a subgraph of *G*, as claimed.

Graph Theory

**Theorem 15.13.** Every simple graph G with  $m \ge 2^{k-3}n$  (that is,  $e(G) \ge 2^{k-2}v(G)$ ) has a  $K_k$ -minor.

**Proof.** We give an inductive proof on the number of edges m of G. If k = 1 then G trivially has a  $K_1$ -minor. If k = 2 and  $m \ge n/2$  (where  $n \ge 2$ , since G is simple) then G trivially has a  $K_2$ -minor. When k = 3 and  $m \ge n$  we have by Theorem 4.3 then G is not a tree and so it contains a cycle. By a sequence of vertex deletions and edge contradictions, we have that G has a  $K_2$ -minor. This we may assume  $k \ge 4$ .

Graph Theory

**Theorem 15.13.** Every simple graph G with  $m \ge 2^{k-3}n$  (that is,  $e(G) \ge 2^{k-2}v(G)$ ) has a  $K_k$ -minor.

**Proof.** We give an inductive proof on the number of edges *m* of *G*. If k = 1 then G trivially has a K<sub>1</sub>-minor. If k = 2 and  $m \ge n/2$  (where n > 2, since G is simple) then G trivially has a  $K_2$ -minor. When k = 3 and m > n we have by Theorem 4.3 then G is not a tree and so it contains a cycle. By a sequence of vertex deletions and edge contradictions, we have that G has a  $K_2$ -minor. This we may assume k > 4. In these cases  $m \ge 2n$ . With n = 1 we have  $m \ge 2n = 2$ , with n = 2 we have m > 2n = 4, with n = 3 we have m > 2n = 6, with n = 4 we have m > 2n = 8, and in each of these cases the result holds vacuously. With n = 5 and  $m \ge 2n = 10$  we must have m = 10 and  $G = K_5$  so that the result holds trivially. Then we may assume that  $k \ge 4$  and  $m \ge 10$ . For the base case, we have that the result holds for all k < 4 and m < 10.

**Theorem 15.13.** Every simple graph G with  $m \ge 2^{k-3}n$  (that is,  $e(G) \ge 2^{k-2}v(G)$ ) has a  $K_k$ -minor.

**Proof.** We give an inductive proof on the number of edges m of G. If k = 1 then G trivially has a K<sub>1</sub>-minor. If k = 2 and  $m \ge n/2$  (where n > 2, since G is simple) then G trivially has a  $K_2$ -minor. When k = 3 and m > n we have by Theorem 4.3 then G is not a tree and so it contains a cycle. By a sequence of vertex deletions and edge contradictions, we have that G has a  $K_2$ -minor. This we may assume k > 4. In these cases m > 2n. With n = 1 we have m > 2n = 2, with n = 2 we have m > 2n = 4, with n = 3 we have m > 2n = 6, with n = 4 we have m > 2n = 8, and in each of these cases the result holds vacuously. With n = 5 and  $m \ge 2n = 10$  we must have m = 10 and  $G = K_5$  so that the result holds trivially. Then we may assume that  $k \ge 4$  and  $m \ge 10$ . For the base case, we have that the result holds for all k < 4 and  $m \le 10$ .

### Theorem 15.13 (continued 1)

**Theorem 15.13.** Every simple graph G with  $m \ge 2^{k-3}n$  (that is,  $e(G) \ge 2^{k-2}v(G)$ ) has a  $K_k$ -minor.

**Proof (continued).** For the induction hypotheses, suppose the result holds for all  $m \le \ell$  (where  $\ell = 10$ ). Let *G* be a graph with *n* vertices and  $m = \ell + 1$  edges where  $m \ge 2^{k-3}n$ . We consider two cases.

First, suppose G has an edge e which lies in at most  $2^{k-3} - 1$  triangles. Then the underlying simple graph of G/e (notice that G/e itself might have multiple edges) has n-1 vertices. Each triangle of G which contains edge e results in a double edge in G/e, which is replaced by a single edge in the underlying simple graph of G/e. So this underlying simple graph has at least  $m - 2^{k-3} \ge 2^{k-3}n - 2^{k-3} = 2^{k-3}(n-1)$  edges. The underlying simple graph of G/e has less edges than graph G, so by the induction hypothesis it has a  $K_k$ -minor. This  $K_k$ -minor of G/e is also a minor of G and the induction step holds in this case.

## Theorem 15.13 (continued 1)

**Theorem 15.13.** Every simple graph G with  $m \ge 2^{k-3}n$  (that is,  $e(G) \ge 2^{k-2}v(G)$ ) has a  $K_k$ -minor.

**Proof (continued).** For the induction hypotheses, suppose the result holds for all  $m \le \ell$  (where  $\ell = 10$ ). Let G be a graph with n vertices and  $m = \ell + 1$  edges where  $m \ge 2^{k-3}n$ . We consider two cases.

First, suppose G has an edge e which lies in at most  $2^{k-3} - 1$  triangles. Then the underlying simple graph of G/e (notice that G/e itself might have multiple edges) has n-1 vertices. Each triangle of G which contains edge e results in a double edge in G/e, which is replaced by a single edge in the underlying simple graph of G/e. So this underlying simple graph has at least  $m - 2^{k-3} \ge 2^{k-3}n - 2^{k-3} = 2^{k-3}(n-1)$  edges. The underlying simple graph of G/e has less edges than graph G, so by the induction hypothesis it has a  $K_k$ -minor. This  $K_k$ -minor of G/e is also a minor of G and the induction step holds in this case.

# Theorem 15.13 (continued 2)

**Proof (continued).** Second suppose that each edge of G lies in at least  $2^{k-3}$  triangles. For  $e \in E$ , denote by t(e) the number of triangles containing e. Any edge e of G is in the subgraph G[N(v)] induced by the neighbors of a vertex v if and only if v is the 'apex' of a triangle whose 'base' is e (i.e., v is a vertex of some triangle containing e, other than the two ends of edge e). Notice that edge e could be in several subgraphs G[N(v)]; in fact, it is such a subgraph t(e) times (i.e., for t(e) different vertices v). Hence we have

$$\sum_{v \in V} |E(G[N(v)])| = \sum_{e \in E} t(e) = 2^{k-3}m \text{ since each of the } m \text{ edges}$$

lies in at least  $2^{k-3}$  triangles in this case

$$= 2^{k-3} \sum_{v \in V} d(v)/2 \text{ by Theorem 1.1}$$
$$= \sum 2^{k-4} d(v).$$

# Theorem 15.13 (continued 2)

**Proof (continued).** Second suppose that each edge of G lies in at least  $2^{k-3}$  triangles. For  $e \in E$ , denote by t(e) the number of triangles containing e. Any edge e of G is in the subgraph G[N(v)] induced by the neighbors of a vertex v if and only if v is the 'apex' of a triangle whose 'base' is e (i.e., v is a vertex of some triangle containing e, other than the two ends of edge e). Notice that edge e could be in several subgraphs G[N(v)]; in fact, it is such a subgraph t(e) times (i.e., for t(e) different vertices v). Hence we have

$$\sum_{v \in V} |E(G[N(v)])| = \sum_{e \in E} t(e) = 2^{k-3}m \text{ since each of the } m \text{ edges}$$
  
lies in at least  $2^{k-3}$  triangles in this case  
$$= 2^{k-3} \sum_{v \in V} d(v)/2 \text{ by Theorem 1.1}$$
  
$$= \sum_{v \in V} 2^{k-4} d(v).$$

# Theorem 15.13 (continued 3)

**Theorem 15.13.** Every simple graph G with  $m \ge 2^{k-3}n$  (that is,  $e(G) \ge 2^{k-2}v(G)$ ) has a  $K_k$ -minor.

**Proof (continued).** This inequality implies that G has at least one vertex v such that its neighborhood subgraph H = G[N(v)] satisfies

$$e(H) \ge 2^{k-4}d(v) = 2^{k-4}v(H).$$
 (\*)

Since *H* is a subgraph of *G* which excludes vertex *v* and all edges incident to *v*, then *H* has less edges than *G*. So, by the induction hypothesis (and (\*)), *H* has a  $K_{k-1}$ -minor. This  $K_{k-1}$  along with vertex *v* and all edges between *v* and vertices of  $K_{k-1}$  (i.e.,  $K_{k-1} \vee K_1$  where  $V(K_1) = \{v\}$ ) gives a  $K_k$ -minor of *G*. In both cases, *G* has a  $K_k$ -minor and so by mathematical induction the claim holds for all graphs satisfying the hypotheses.

# **Corollary 15.14.** For $k \ge 2$ , every $(2^{k-2} + 1)$ -chromatic graph has a $K_k$ -minor.

**Proof.** Let *G* be a  $(2^{k-2} + 1)$ -chromatic graph and let *F* be a  $(2^{k-2} + 1)$ -critical subgraph of *G*. By Theorem 14.7,  $\delta(F) \ge (2^{k-2} + 1) - 1 = 2^{k-2}$ , and so  $e(F) \ge \delta(F)v(F)/2 \ge 2^{2k-3}v(F)$ . Hence by Theorem 15.7, *F* has a  $K_k$ -minor. Therefore, *G* has a  $K_k$ -minor, as claimed.

**Corollary 15.14.** For  $k \ge 2$ , every  $(2^{k-2} + 1)$ -chromatic graph has a  $K_k$ -minor.

**Proof.** Let *G* be a  $(2^{k-2} + 1)$ -chromatic graph and let *F* be a  $(2^{k-2} + 1)$ -critical subgraph of *G*. By Theorem 14.7,  $\delta(F) \ge (2^{k-2} + 1) - 1 = 2^{k-2}$ , and so  $e(F) \ge \delta(F)v(F)/2 \ge 2^{2k-3}v(F)$ . Hence by Theorem 15.7, *F* has a  $K_k$ -minor. Therefore, *G* has a  $K_k$ -minor, as claimed.

()