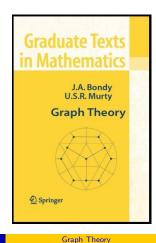
## **Graph Theory**

### Chapter 16. Matchings

16.2. Matchings in Bipartite Graphs—Proofs of Theorems



Theorem 16.4. Hall's Theorem

## Theorem 16.4 (continued 1)

**Proof (continued).** Since  $u \in X$  is not covered by  $M^*$ , for any  $z \in Z$  with  $z \neq u$  we have that there is an  $M^*$ -alternating path from u to z (which starts with an edge NOT in  $M^*$ ). If z is not covered by  $M^*$ , then the  $M^*$ -alternating path from u to z is in fact an  $M^*$ -augmented path in G, contradicting the fact that G contains no  $M^*$ -augmented paths. Hence, z is covered by  $M^*$  and so u is the only element of Z not covered by  $M^*$ .

Define  $R = X \cap Z$  and  $B = Y \cap Z$  (see Figure 16.6).

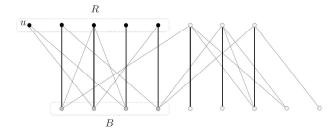


Fig. 16.6. Proof of Hall's Theorem (16.4)

### Theorem 16.4. Hall's Theorem

#### Theorem 16.4. Hall's Theorem.

A bipartite graph G = G[X, Y] has a matching which covers every vertex in X if and only if  $|N(S)| \ge |S|$  for all  $S \subseteq X$  (where N(S) is the set of all vertices which are neighbors of some vertex in S).

**Proof.** Let G = G[X, Y] be a bipartite graph which has a matching M covering every vertex in X. Let  $S \subseteq X$ . The vertices in S are matched in M with distinct vertices in N(S), defining an injective ("one to one") function from S to N(S). So  $|N(S)| \ge |S|$ , as claimed.

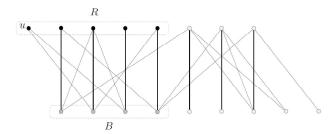
Conversely, let G = G[X, Y] be a bipartite graph which has no matching covering every vertex in X. Let  $M^*$  be a maximum matching in G and let u be a vertex in X not covered by  $M^*$ . Let Z denote the set of all vertices reachable from u by  $M^*$ -alternating paths. Because  $M^*$  is a maximum matching, by Theorem 16.3 G contains no  $M^*$ -augmenting path.

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#### Theorem 16.4. Hall's Theorer

# Theorem 16.4 (continued 2)

### Proof (continued).



**Fig. 16.6.** Proof of Hall's Theorem (16.4)

Now the vertices of  $R \setminus \{u\}$  are matched under  $M^*$  with the vertices of B (because of the  $M^*$ -alternating path definition of Z). This implies a bijection between  $R \setminus \{u\}$  and B so that |B| = |R| - 1. Now the neighbors of vertices in R include all vertices in B; that is  $N(R) \supset B$ . In fact, every vertex in  $N(R) \subset Z$  is connected to u by an  $M^*$ -alternating path, and so N(R) = B. Hence |N(R)| = |B| = |R| - 1.

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# Theorem 16.4 (continued 3)

#### Theorem 16.4. Hall's Theorem.

A bipartite graph G = G[X, Y] has a matching which covers every vertex in X if and only if  $|N(S)| \ge |S|$  for all  $S \subseteq X$  (where N(S) is the set of all vertices which are neighbors of some vertex in S).

**Proof (continued).** So with set S equal to set R, we have |N(R)| = |N(S)| < |S| = |R|. That is, if G = G[X, Y] does not have a matching which covers every vertex in X then |N(S)| < |S| for some  $S \subseteq X$ , as claimed.

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Theorem 8.33

### Theorem 8.32

### Theorem 8.32. The König-Egerváry Theorem.

In any bipartite graph G, the number of edges in a maximum matching is equal to the number of vertices in a minimum covering.

**Proof.** Let G = G[X, Y] be a bipartite graph with  $M^*$  a maximum matching in G, and U the set of vertices in X not covered by  $M^*$ . Denote by Z the set of all vertices in G reachable from some vertex in G by G by G by G by G alternating paths.

Define  $R = X \cap Z$  and  $B = Y \cap Z$ . Let  $K^* = (Z \setminus R) \cup B$ . Then  $K^*$  is a covering with  $|K^*| = |M^*|$  by Exercise 16.2.8. By Proposition 16.7,  $K^*$  is a minimum cover. That is,  $\alpha'(G) = |M^*| = |K^*| = \beta(G)$ , as claimed.  $\square$ 

## Corollary 16.6

**Corollary 16.6.** Every nonempty regular bipartite graph has a perfect matching.

**Proof.** First, if G[X,Y] is k-regular where  $k \geq 1$  then |X| = |Y| by Exercise 1.1.9. Let  $S \subseteq X$  and let  $E_1$  and  $E_2$  denote the sets of edges of G incident with S and G(S), respectively. Notice that G(S) (but not edge with one end in G(S) must have the other end in G(S) (but not conversely). Since G(S) is G(S) is G(S) is arbitrary, by Corollary 16.5 we have that G(S) has a perfect matching.

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