Graph Theory

Chapter 18. Hamilton Cycles

18.1. Hamiltonian and Nonhamiltonian Graphs-Proofs of Theorems



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Theorem 18.1

Theorem 18.1. Let *S* be a set of vertices of a hamiltonian graph *G*. Then the number of connected component of G - S satisfies $c(G - S) \le |S|$. Moreover, if equality holds, then each of the |S| components of G - S is traceable, and every Hamilton cycle of *G* includes a Hamilton path in each of these components.

Proof. Let *C* be a Hamilton cycle of *G*. Then C - S has at most |S| connected components. Since *C* is a spanning subgraph of *G* then G - S also has at most |S| components. That is, $c(G - S) \le |S|$, as claimed.

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If equality holds and G - S has exactly |S| components, then C - S also has exactly |S| components (since this is the maximum number of components and G - S has more edges than C - S). The connected components of C - S are then spanning subgraphs of the connected components of G - S.

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Proof (continued). Since the components of C - S are paths (in fact, spanning paths of the components of C - S) then the components of G - S include Hamilton paths consisting of components of C - S. Therefore, the components of G - S are traceable, as claimed.