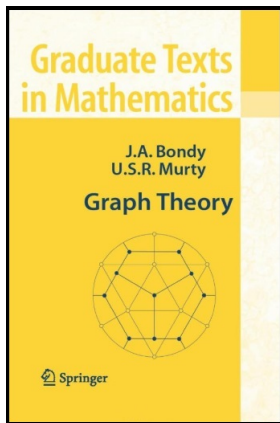


# Graph Theory

## Chapter 2. Subgraphs

### 2.1. Subgraphs and Supergraphs—Proofs of Theorems



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# Theorem 2.1

**Theorem 2.1.** Let  $G$  be a graph in which all vertices have degree at least two. Then  $G$  contains a cycle.

**Proof.** If  $G$  is not simple then it either contains a loop (i.e., a cycle of length one) or parallel edges (two of which with the same ends form a cycle of length two). So we can assume without loss of generality that  $G$  is simple.

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Let  $P = v_0 v_1 \cdots v_{k-1} v_k$  be a path in  $G$  of longest length (which exists since  $G$  is finite). Since vertex  $v_k$  is of degree at least two by hypothesis then it has a neighbor  $v$  different from  $v_{k-1}$ . If  $v$  is not on  $P$ , then the path  $v_0 v_1 \cdots v_k v$  is larger than path  $P$ , contradicting the choice of  $P$ . So it must be that  $v = v_i$  for some  $0 \leq i \leq k - 2$ . Then  $G$  contains the cycle  $v_i v_{i+1} \cdots v_k v_i$ .  $\square$

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## Theorem 2.2

**Theorem 2.2.** Any simple graph  $G$  with  $\sum_{v \in V} \binom{d(v)}{2} > \binom{n}{2}$  contains a quadrilateral.

**Proof.** Denote by  $p_2$  the number of distinct paths of length 2 in  $G$ . Denote by  $p_2(v)$  the number of such paths whose “central” vertex is  $v$ . Now for a given vertex  $v$  where  $d(v) \geq 2$ , we can choose 2-paths with  $v$  as the central vertex in  $\binom{d(v)}{2}$  different ways.

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So for each vertex  $v$  we have  $p_2(v) = \binom{d(v)}{2}$  (where we interpret  $\binom{0}{2} = \binom{1}{2} = 0$ ). Since each 2-path has a unique central vertex, then

$$p_2 = \sum_{v \in V} p_2(v) = \sum_{v \in V} \binom{d(v)}{2}. \quad (*)$$

Next, each 2-path has a unique pair of end vertices. So we can create sets of 2-paths where two 2-paths are in the same set if they have the same end vertices.

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## Theorem 2.2 (continued)

**Theorem 2.2.** Any simple graph  $G$  with  $\sum_{v \in V} \binom{d(v)}{2} > \binom{n}{2}$  contains a quadrilateral.

**Proof (continued).** There are then  $\binom{n}{2}$  such sets (though some could be empty). We have hypothesized  $\sum_{v \in V} \binom{d(v)}{2} > \binom{n}{2}$ . So by (\*),  $p_2 > \binom{n}{2}$  and by the Pigeonhole Principle one of the sets of 2-paths with common end vertices must contain at least two 2-paths. Since these two 2-paths have different central vertices (because they are elements of a set) then the union of these two 2-paths forms a quadrilateral that is contained in  $G$ , as claimed.  $\square$