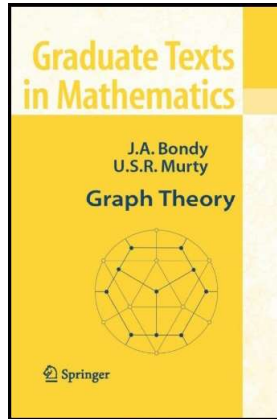


Graph Theory

Chapter 2. Subgraphs

2.5. Edge Cuts and Bonds—Proofs of Theorems



Theorem 2.10

Theorem 2.10. A graph $G = (V, E)$ is even if and only if $|\partial(X)|$ is even for every subset X of V .

Proof. Suppose $|\partial(X)|$ is even for every subset X of V . Then $|\partial(v)|$ is even for every vertex v (here, we take $X = \{v\}$). In a loopless graph, $|\partial(v)| = d(v)$ since $\partial(v)$ is the set of all links incident to v . Since each loop with v as the end contributes 2 to the degree, then $d(v)$ is even if G has loops. So every vertex of G is of even degree and hence G is even.

Suppose G is even. Then (by definition of “even”) $d(v)$ is even for each $v \in V$, and so $\sum_{v \in V} d(v)$ is even. So by Theorem 2.9, it follows that $|\partial(X)|$ must be even for every $X \subset V$. \square

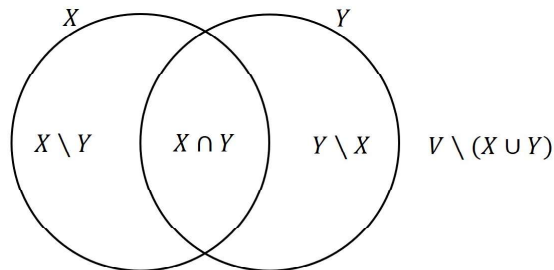
Proposition 2.11

Proposition 2.11

Proposition 2.11. Let G be a graph and let X and Y be subsets of V . Then $\partial(X) \Delta \partial(Y) = \partial(X \Delta Y)$.

Proof. Bondy and Murty give crude Venn diagrams to explain this proof. We also give a “proof by picture,” but we use color coded edges to describe edge cuts. First, for given $X, Y \subset V$ we partition V as

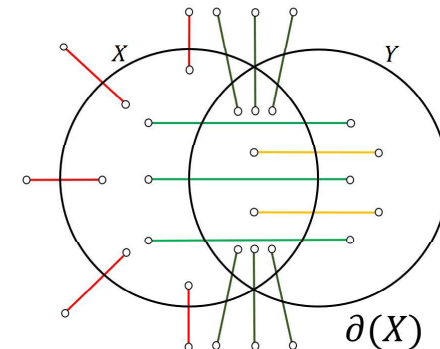
$$V = (X \setminus Y) \cup (Y \setminus X) \cup (X \cap Y) \cup (V \setminus (X \cup Y)) :$$



Proposition 2.11

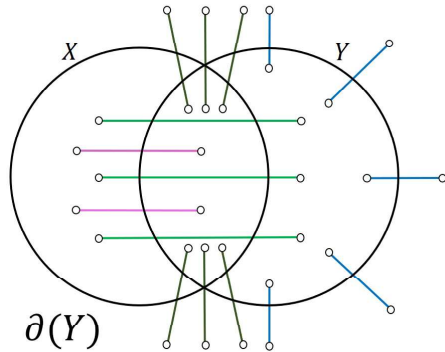
Proposition 2.11 (continued 1)

Proof (continued). Based on these four disjoint sets, we now describe $\partial(X)$, $\partial(Y)$, and $\partial(X \Delta Y)$. The edges in $\partial(X)$ are of the following four types: (1) edges with one end in $X \setminus Y$ and one end in $V \setminus (X \cup Y)$ (in red), (2) edges with one end in $X \setminus Y$ and one end in $Y \setminus X$ (in green), (3) edges with one end in $X \cap Y$ and one end in $V \setminus (X \cup Y)$ (in brown), and (4) edges with one end in $X \cap Y$ and one end in $Y \setminus X$ (in orange).



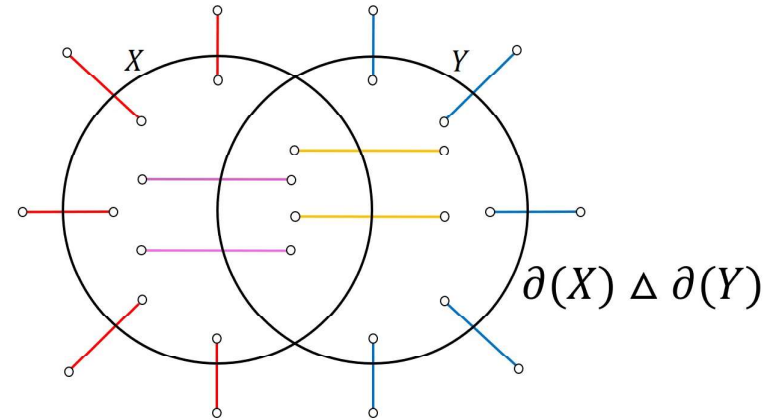
Proposition 2.11 (continued 2)

Proof (continued). The edges in $\partial(Y)$ are of the following four types: (5) edges with one end in $Y \setminus X$ and one end in $V \setminus (X \cup Y)$ (in blue), (6) edges with one end in $Y \setminus X$ and one end in $X \setminus Y$ (in green), (7) edges with one end in $X \cap Y$ and one end in $V \setminus (X \cup Y)$ (in brown), and (8) edges with one end in $X \cap Y$ and one end in $X \setminus Y$ (in pink).



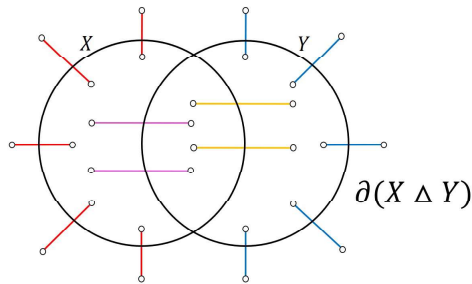
Proposition 2.11 (continued 3)

Proof (continued). Now $\partial(X) \Delta \partial(Y) = (\partial(X) \setminus \partial(Y)) \cup (\partial(Y) \setminus \partial(X))$ consists of the edges colored red, orange, blue, and pink.



Proposition 2.11 (continued 4)

Proof (continued). Next, $X \Delta Y = (X \setminus Y) \cup (Y \setminus X)$ and the edges of $\partial(X \Delta Y)$ are of the following four types: (9) edges with one end in $X \setminus Y$ and one end in $V \setminus (X \cup Y)$ (in red), (10) edges with one end in $X \setminus Y$ and one end in $X \cap Y$ (in pink), (11) edges with one end in $Y \setminus X$ and one end in $V \setminus (X \cup Y)$ (in blue), and (12) edges with one end in $Y \setminus X$ and one end in $X \cap Y$ (in orange).



So $\partial(X \Delta Y)$ consists of the edges colored red, orange, blue, and pink. Therefore, $\partial(X) \Delta \partial(Y) = \partial(X \Delta Y)$, as claimed. \square

Theorem 2.15

Theorem 2.15. In a connected graph G , a nonempty edge cut $\partial(X)$ is a bond if and only if both $G[X]$ and $G[V \setminus X]$ are connected.

Proof. First, suppose $\partial(X)$ is a bond. Let Y be a nonempty proper subset of X . Since G is connected by hypothesis, then both $\partial(Y)$ and $\partial(X \setminus Y)$ are nonempty (or else Y and $V \setminus Y$, or $X \setminus Y$ and $V \setminus (X \setminus Y)$ form a "separation" of G , contradicting the fact that G is connected). If $E[Y, X \setminus Y]$ is empty (i.e., if there are no edges of G with one end in Y and one end in $X \setminus Y$) then $\partial(Y)$ is a proper nonempty subset of $\partial(X)$ (since in this case $\partial(Y)$ consists only of edges with one end in Y and one end in $X \setminus Y$, and with $Y \subseteq X$ all these edges are in $\partial(X)$). So we must have $E[Y, X \setminus Y] \neq \emptyset$. Since Y is an arbitrary nonempty proper subset of X , then there is no "separation" of the induced subgraph $G[X]$ of G and hence $G[X]$ is connected. Since $\partial(X) = \partial(V \setminus X)$ and since $\partial(X)$ is a bond then $\partial(V \setminus X)$ is a bond. So a similar argument (based on Y a proper nonempty subset of $V \setminus X$) shows that $G[V \setminus X]$ is connected.

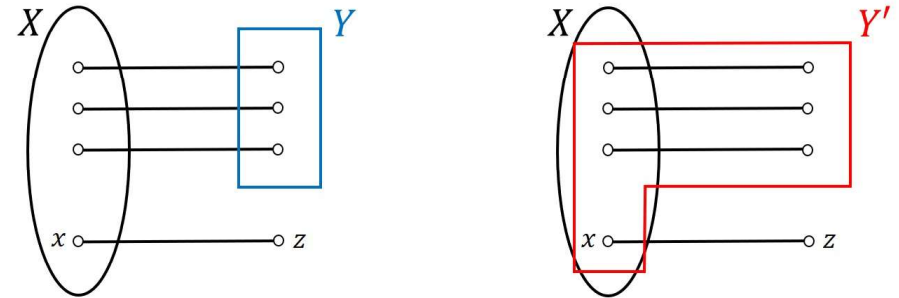
Theorem 2.15 (continued 1)

Proof (continued). Second, suppose $\partial(X)$ is not a bond. Then $\partial(X)$ is not a minimal edge cut. So there is some subset Y of V with $\emptyset \neq \partial(Y) \subsetneq \partial(X)$, $Y \neq \emptyset$ (since $\partial(\emptyset) = \emptyset$), and $Y \neq V$ (since $\partial(V) = \emptyset$).

We claim, without loss of generality, that $X \cap Y \neq \emptyset$. Suppose, to the contrary, that $X \cap Y = \emptyset$. Then if $e \in \partial(Y) \subseteq \partial(X)$ then one end of e is in Y and one end is in $V \setminus Y$, AND one end of e is in X and one end is in $V \setminus X$. In the event that $X \cap Y = \emptyset$ then e cannot have one end in X and one end in $V \setminus (X \cup Y)$ (since then $e \notin \partial(Y)$) nor can e have one end in Y and one end in $V \setminus (X \cup Y)$ (since then $e \notin \partial(X)$). So for $e \in \partial(Y)$, e has one end in X and one end in Y ; in fact, $\partial(Y) = E[X, Y]$. If $e \in \partial(X) \setminus \partial(Y)$ then e has one end in X and one end in $V \setminus (X \cup Y)$. So the fact that $\partial(Y) \subsetneq \partial(X)$ implies that there is an edge in $\partial(X)$ of the form $\{x, z\}$ where $x \in X$ and $z \in V \setminus (X \cup Y)$.

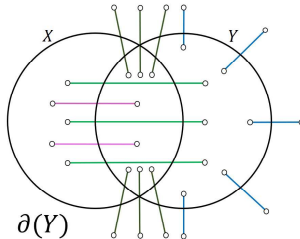
Theorem 2.15 (continued 2)

Proof (continued). Let $Y' = X \cup Y$. Then $\partial(Y') = \partial(X) \setminus E[X, Y] \subsetneq \partial(X)$ (since $E[X, Y] = \partial(Y) \neq \emptyset$ here). Since edge $\{x, z\} \in \partial(X \cup Y) = \partial(Y')$ then $Y' \neq V$. So we can replace Y with Y' where $X \cap Y' \neq \emptyset$. So we may assume without loss of generality that $X \cap Y \neq \emptyset$ above, as claimed.



Theorem 2.15 (continued 3)

Proof (continued). Since $\partial(Y) \subset \partial(X)$ then we must have $E[X \cap Y, X \setminus Y] = \emptyset$ (in pink) and $E[Y \setminus X, V \setminus (X \cup Y)] = \emptyset$ (in blue).



Now if $X \setminus Y \neq \emptyset$ then $X \cap Y$ and $X \setminus Y$ form a “separation” of $G[X]$ and so $G[X]$ is not connected (this is where we need $X \cap Y \neq \emptyset$). If $X \setminus Y = \emptyset$ (in which case $X \subsetneq Y$) then $\emptyset \neq Y \setminus X$ (since $\partial(Y) \subsetneq \partial(X)$ then $X \neq Y$) and $\emptyset \neq Y \setminus X \subsetneq V \setminus X$, so that X and $Y \setminus X$ form a “separation” of $G[V \setminus X]$ and $G[V \setminus X]$ is not connected. That is, if $\partial(X)$ is not a bond then either $G[X]$ is not connected or $G[V \setminus X]$ is not connected. \square